Advanced CUDA: Application Examples

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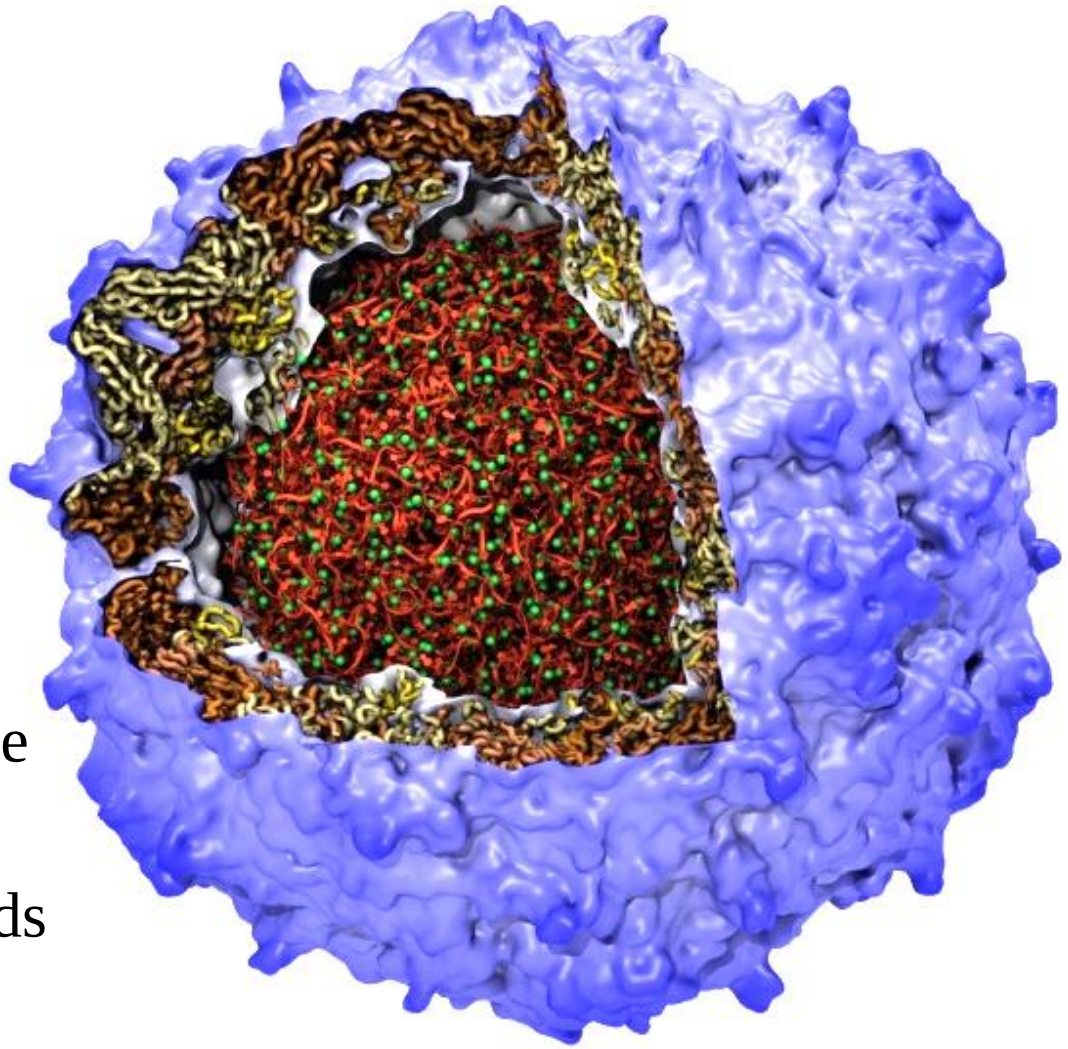
http://www.ks.uiuc.edu/Research/gpu/

GPGPU2: Advanced Methods for Computing with CUDA,

University of Cape Town, April 2014

Molecular Surface Visualization

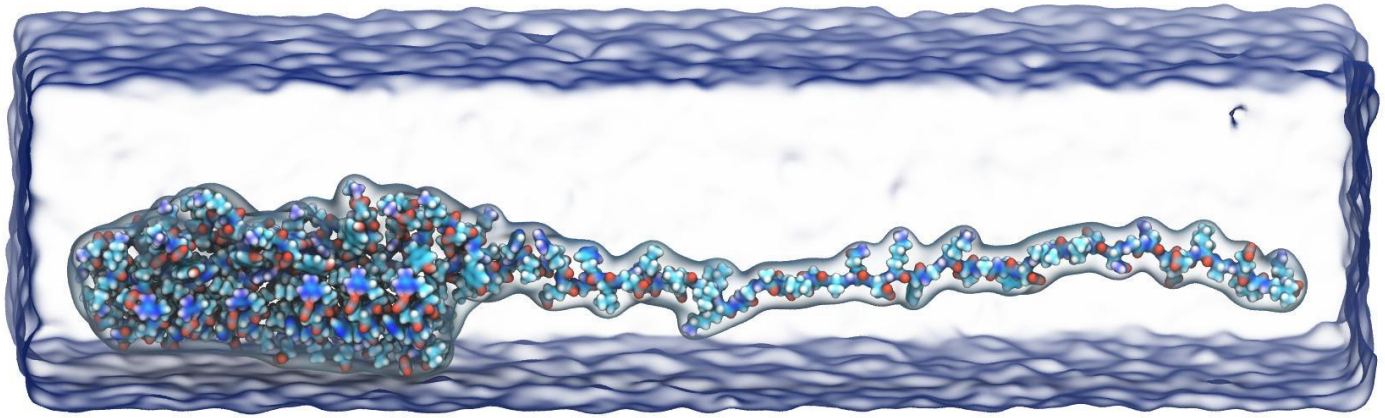
- Large biomolecular complexes are difficult to interpret with atomic detail graphical representations
- Even secondary structure representations become cluttered
- Surface representations are easier to use when greater abstraction is desired, but are computationally costly
- Most surface display methods incapable of animating dynamics of large structures w/ millions of particles



Poliovirus

VMD "QuickSurf" Representation

- Displays continuum of structural detail:
  - All-atom models
  - Coarse-grained models
  - Cellular scale models
  - Multi-scale models: All-atom + CG, Brownian + Whole Cell
  - Smoothly variable between full detail, and reduced resolution representations of very large complexes

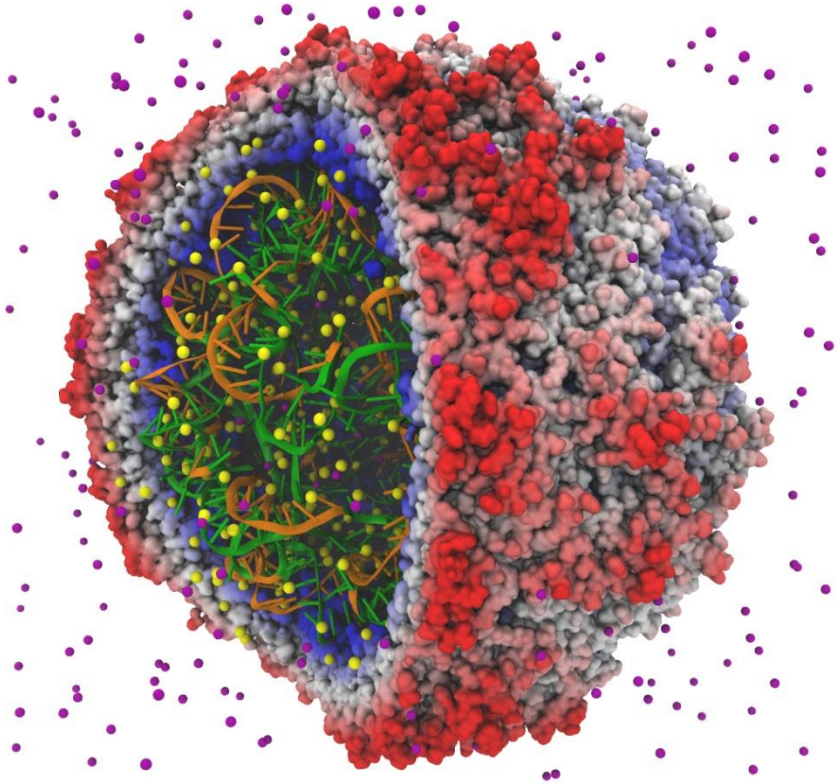


Fast Visualization of Gaussian Density Surfaces for Molecular Dynamics and Particle System Trajectories.

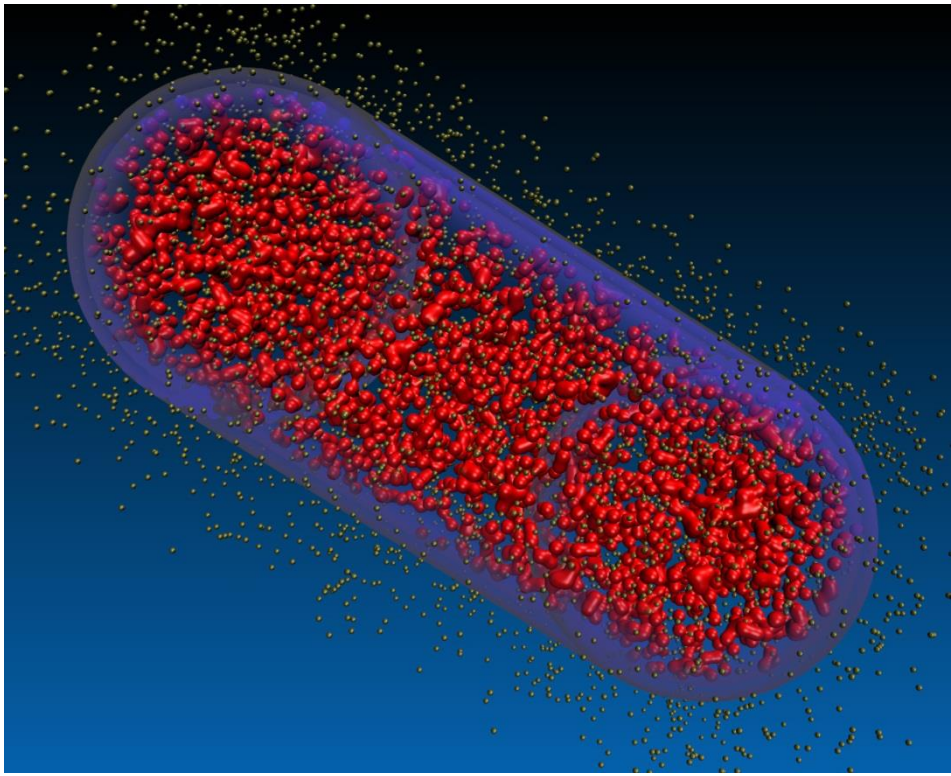
M. Krone, J. E. Stone, T. Ertl, K. Schulten. *EuroVis Short Papers*, pp. 67-71, 2012

VMD “QuickSurf” Representation

- Uses multi-core CPUs and GPU acceleration to enable **smooth real-time animation** of MD trajectories
- Linear-time algorithm, scales to millions of particles, as limited by memory capacity

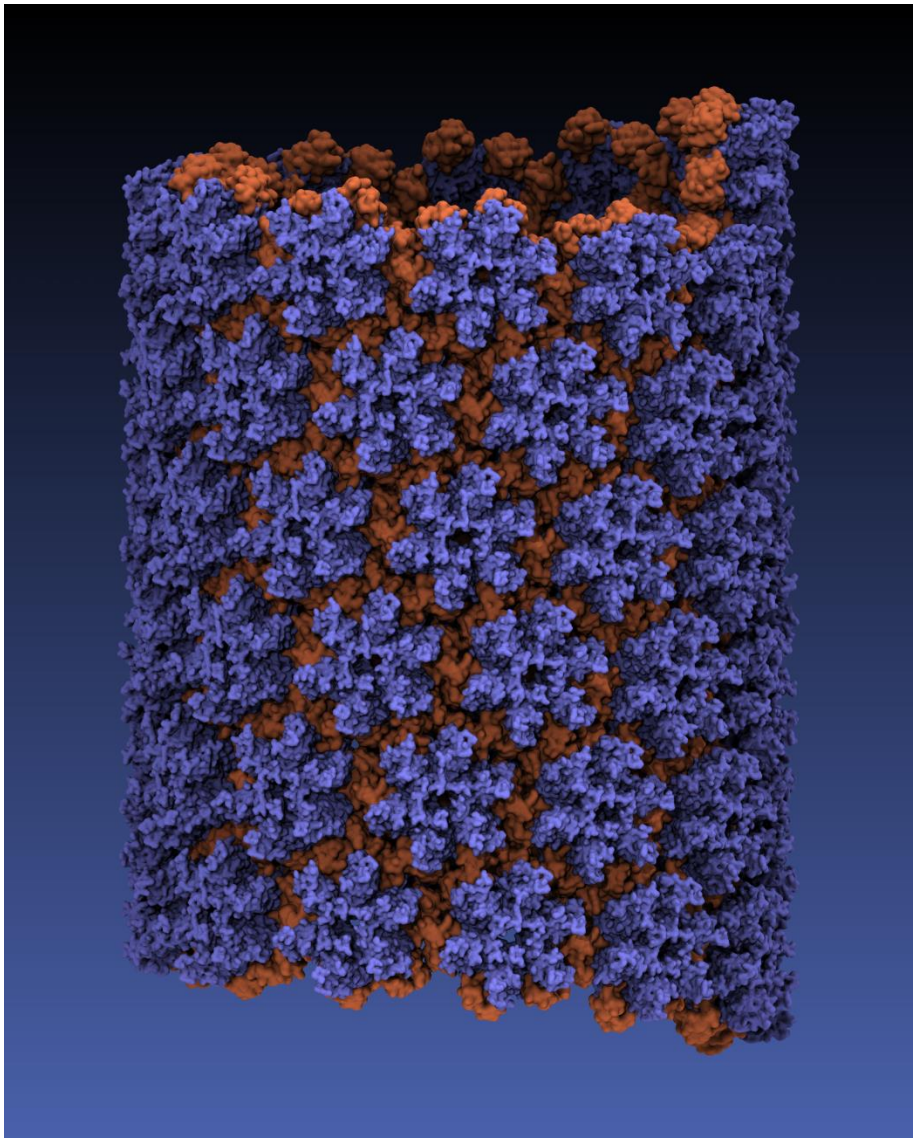


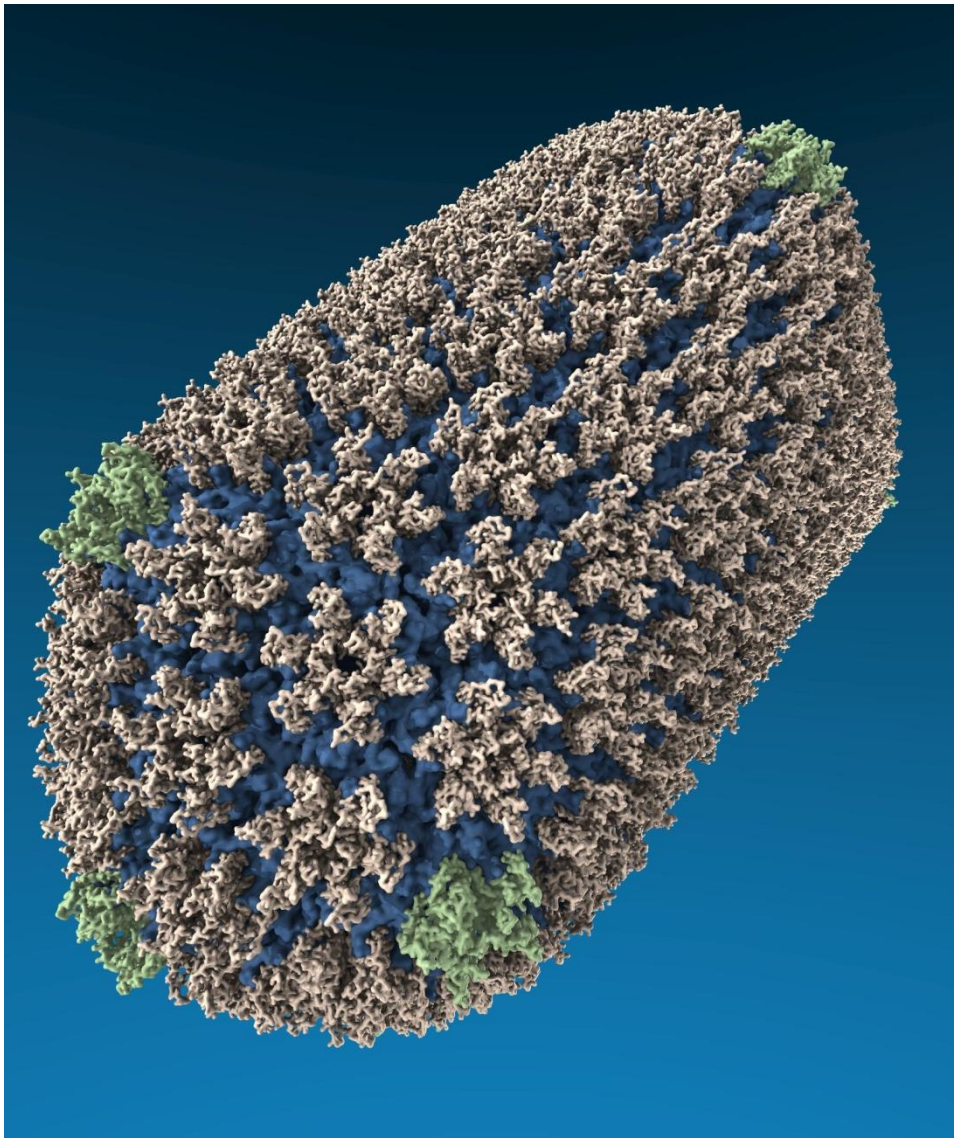
Satellite Tobacco Mosaic Virus



Lattice Cell Simulations

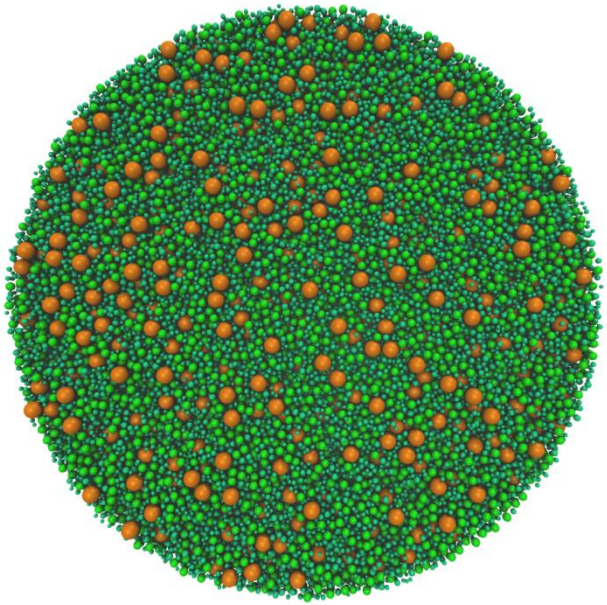
VMD “QuickSurf” Representation

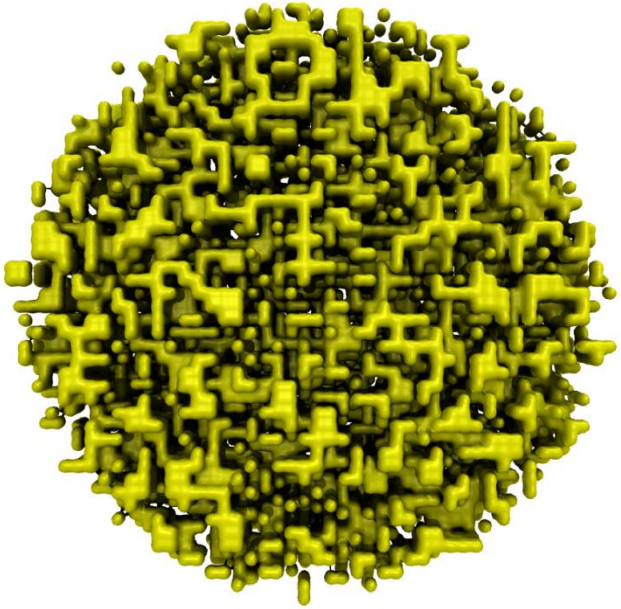


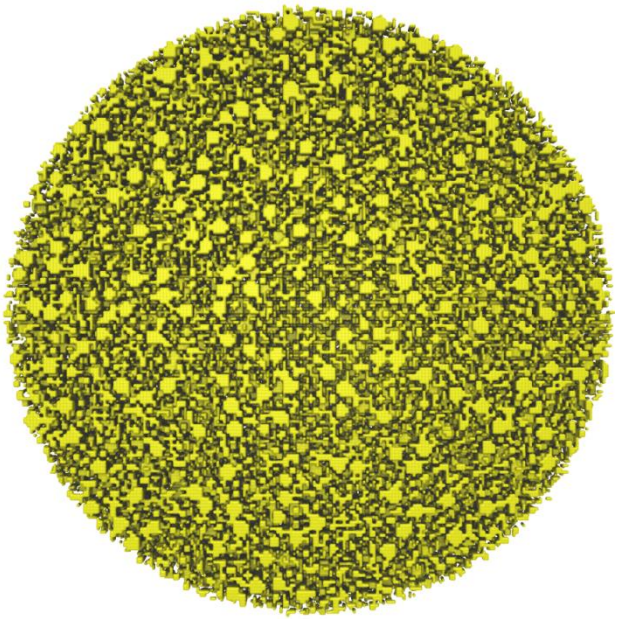


All-atom HIV capsid simulations

QuickSurf Representation of Lattice Cell Models







Continuous particle based model – often 70 to 300 million particles

Discretized lattice models derived from continuous model shown in VMD QuickSurf representation

Lattice Microbes: High-performance stochastic simulation method for the reaction-diffusion master equation

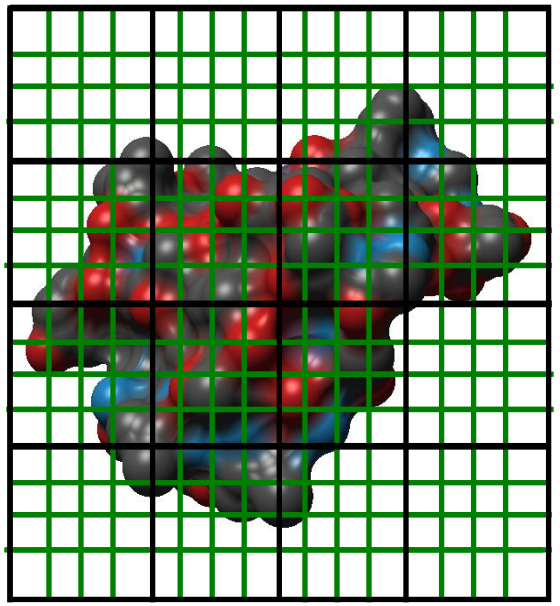
E. Roberts, J. E. Stone, and Z. Luthey-Schulten.
J. Computational Chemistry 34 (3), 245-255, 2013.

QuickSurf Algorithm Overview

- Build spatial acceleration data structures, optimize data for GPU
- Compute 3-D density map, 3-D volumetric texture map:

$$\rho(\vec{r}; \vec{r}_1, \vec{r}_2, \ldots, \vec{r}_N) = \sum_{i=1}^{N} e^{\frac{-|\vec{r}-\vec{r}_i|^2}{2\alpha^2}}$$

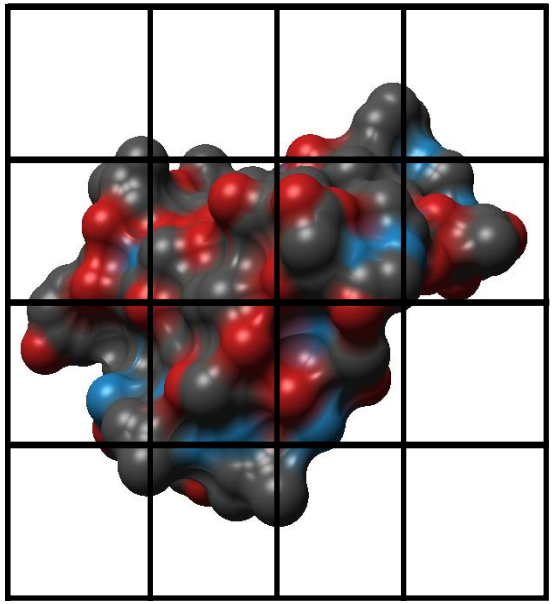
- Extract isosurface for a user-defined density value



3-D density map lattice, spatial acceleration grid, and extracted surface

QuickSurf Particle Sorting, Bead Generation, Spatial Hashing

- Particles sorted into spatial acceleration grid:
  - Selected atoms or residue “beads” converted lattice coordinate system
  - Each particle/bead assigned cell index, sorted w/NVIDIA Thrust template library
- Complication:
  - Thrust allocates GPU mem. on-demand, no recourse if insufficient memory, have to re-gen QuickSurf data structures if caught by surprise!
- Workaround:
  - Pre-allocate guesstimate workspace for Thrust
  - Free the Thrust workspace right before use
  - Newest Thrust allows user-defined allocator code…



Coarse resolution spatial acceleration grid

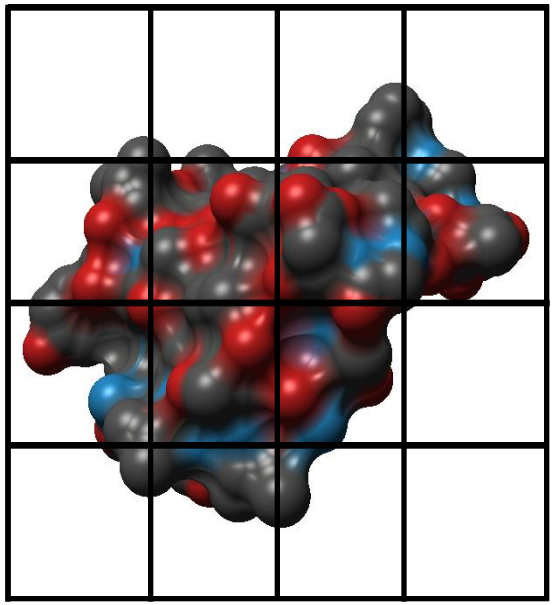
Spatial Hashing Algorithm Steps/Kernels

1) Compute bin index for each atom, store to memory w/ atom index

2) **Sort** list of bin and atom index tuples (1) by bin index **(thrust kernel)**

3) Count atoms in each bin (2) using a **parallel prefix sum, aka *scan***, compute the destination index for each atom, store per-bin starting index and atom count **(thrust kernel)**

4) Write atoms to the output indices computed in (3), and we have completed the data structure



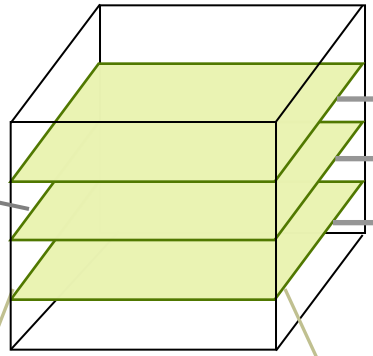
QuickSurf uniform grid spatial subdivision data structure

QuickSurf and Limited GPU Global Memory

- High resolution molecular surfaces require a fine lattice spacing
- Memory use grows cubically with decreased lattice spacing
- Not typically possible to compute a surface in a single pass, so we loop over sub-volume “chunks” until done…
- Chunks pre-allocated and sized to GPU global mem capacity to prevent unexpected memory allocation failure while animating…
- Complication:
  - Thrust allocates GPU mem. on-demand, no recourse if insufficient memory, have to re-gen QuickSurf data structures if caught by surprise!
- Workaround:
  - Pre-allocate guesstimate workspace for Thrust
  - Free the Thrust workspace right before use
  - Newest Thrust allows user-defined allocator code…

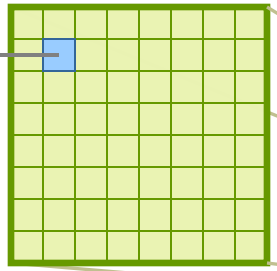
QuickSurf Density Parallel Decomposition

QuickSurf 3-D density map decomposes into thinner 3-D slabs/slices (CUDA grids)



Large volume computed in multiple passes, or multiple GPUs

Small 8x8 thread blocks afford large per-thread register count, shared memory



Each thread computes one or more density map lattice points

Threads producing results that are used

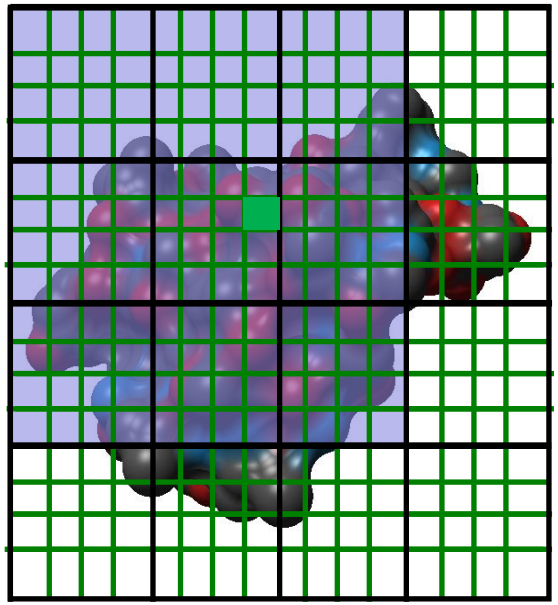
Inactive threads, region of discarded output

Padding optimizes global memory performance, guaranteeing coalesced global memory accesses

Grid of thread blocks

QuickSurf Density Map Algorithm

- Spatial acceleration grid cells are sized to match the cutoff radius for the exponential, beyond which density contributions are negligible
- Density map lattice points computed by summing density contributions from particles in 3x3x3 grid of neighboring spatial acceleration cells
- Volumetric texture map is computed by summing particle colors normalized by their individual density contribution



3-D density map lattice point and the neighboring spatial acceleration cells it references

QuickSurf Density Map Kernel Optimizations

- Compute reciprocals, prefactors, other math on the host CPU prior to kernel launch
- Use of **intN** and **floatN** vector types in CUDA kernels for improved global memory bandwidth
- **Thread coarsening**: one thread computes multiple output densities and colors
- Input data and **register tiling**: share blocks of input, partial distances in regs shared among multiple outputs
- Global memory **(L1 cache) broadcasts**: all threads in the block traverse the same atom/particle at the same time

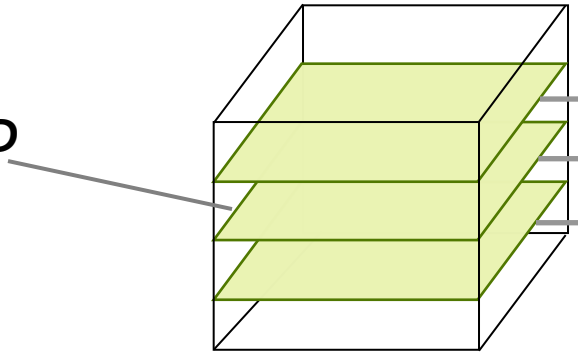
QuickSurf Density Map Kernel Snippet…

```
for (zab=zabmin; zab<=zabmax; zab++) {
  for (yab=yabmin; yab<=yabmax; yab++) {
    for (xab=xabmin; xab<=xabmax; xab++) {
      int abcellidx = zab * acplanesz + yab * acncells.x + xab;
      uint2 atomstartend = cellStartEnd[abcellidx];
      if (atomstartend.x != GRID_CELL_EMPTY) {
        for (unsigned int atomid=atomstartend.x; atomid<atomstartend.y; atomid++) {
          float4 atom = sorted_xyzr[atomid];
          float dx = coorx - atom.x;          float dy = coory - atom.y;        float dz = coorz - atom.z;
          float dxy2 = dx*dx + dy*dy;
          float r21 = (dxy2 + dz*dz) * atom.w;
          densityval1 += exp2f(r21);
          /// Loop unrolling and register tiling benefits begin here……
          float dz2 = dz + gridspacing;
          float r22 = (dxy2 + dz2*dz2) * atom.w;
          densityval2 += exp2f(r22);
          /// More loop unrolling ….
```

QuickSurf Marching Cubes Isosurface Extraction

- Isosurface is extracted from each density map "chunk", and either copied back to the host, or **rendered directly** out of GPU global memory via **CUDA/OpenGL interop**
- All MC memory buffers are pre-allocated to prevent significant overhead when animating a simulation trajectory

QuickSurf 3-D density map decomposes into thinner 3-D slabs/slices (CUDA grids)



…

Chunk 2

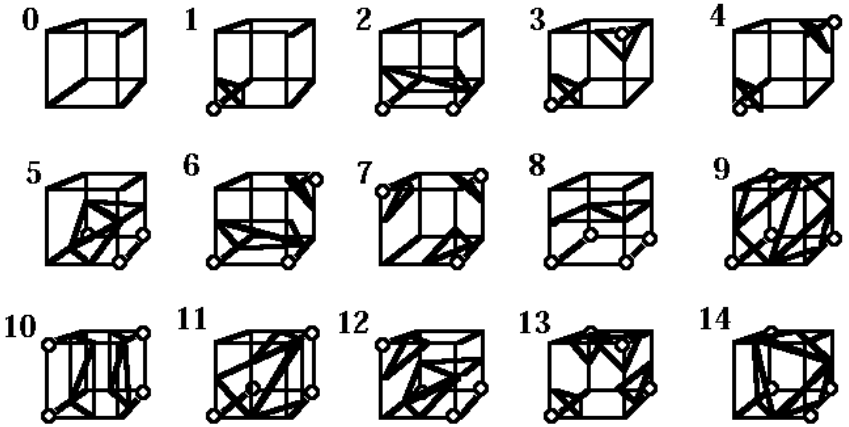
Chunk 1

Chunk 0

Large volume computed in multiple passes

Brief Marching Cubes Isosurface Extraction Overview

- Given a 3-D volume of scalar density values and a requested surface density value, marching cubes computes vertices and triangles that compose the requested surface triangle mesh
- Each MC "cell" (a cube with 8 density values at its vertices) produces a variable number of output vertices depending on how many edges of the cell contain the requested isovalue…
- Use **scan()** to compute the output indices so that each worker thread has **conflict-free output** of vertices/triangles



Brief Marching Cubes Isosurface Extraction Overview

- Once the output vertices have been computed and stored, we compute surface normals and colors for each of the vertices
- Although the separate normals+colors pass reads the density map again, molecular surfaces tend to generate a small percentage of MC cells containing triangles, we avoid wasting interpolation work
- We use CUDA **tex3D()** hardware 3-D texture mapping:
  - Costs double the texture memory and a one copy from GPU global memory to the target texture map with **cudaMemcpy3D()**
  - Still roughly 2x faster than doing color interpolation without the texturing hardware, at least on GT200 and Fermi hardware
  - Kepler has new texture cache memory path that may make it feasible to do our own color interpolation and avoid the use of extra 3-D texture memory and associated copy, with acceptable performance

QuickSurf Marching Cubes Isosurface Extraction

- Our optimized MC implementation computes per-vertex surface normals, colors, and outperforms the NVIDIA SDK sample by a fair margin on Fermi GPUs
- Complications:
  - Even on a 6GB Quadro 7000, GPU global memory is under great strain when working with large molecular complexes, e.g. viruses
  - Marching cubes involves a parallel prefix sum (scan) to compute target indices for writing resulting vertices
  - We use Thrust for scan, has the same memory allocation issue mentioned earlier for the sort, so we use the same workaround
  - The number of output vertices can be huge, but we rarely have sufficient GPU memory for this – we use a fixed size vertex output buffer and hope our heuristics don't fail us

QuickSurf Performance
GeForce GTX 580

Molecular system	Atoms	Resolution	T_{sort}	$T_{density}$	T_{MC}	# vertices	FPS
MscL	111,016	1.0Å	0.005	0.023	0.003	0.7 M	28
STMV capsid	147,976	1.0Å	0.007	0.048	0.009	2.4 M	13.2
Poliovirus capsid	754,200	1.0Å	0.01	0.18	0.05	9.2 M	3.5
STMV w/ water	955,225	1.0Å	0.008	0.189	0.012	2.3 M	4.2
Membrane	2.37 M	2.0Å	0.03	0.17	0.016	5.9 M	3.9
Chromatophore	9.62 M	2.0Å	0.16	0.023	0.06	11.5 M	3.4
Membrane w/ water	22.77 M	4.0Å	4.4	0.68	0.01	1.9 M	0.18

Fast Visualization of Gaussian Density Surfaces for Molecular Dynamics and Particle System Trajectories.

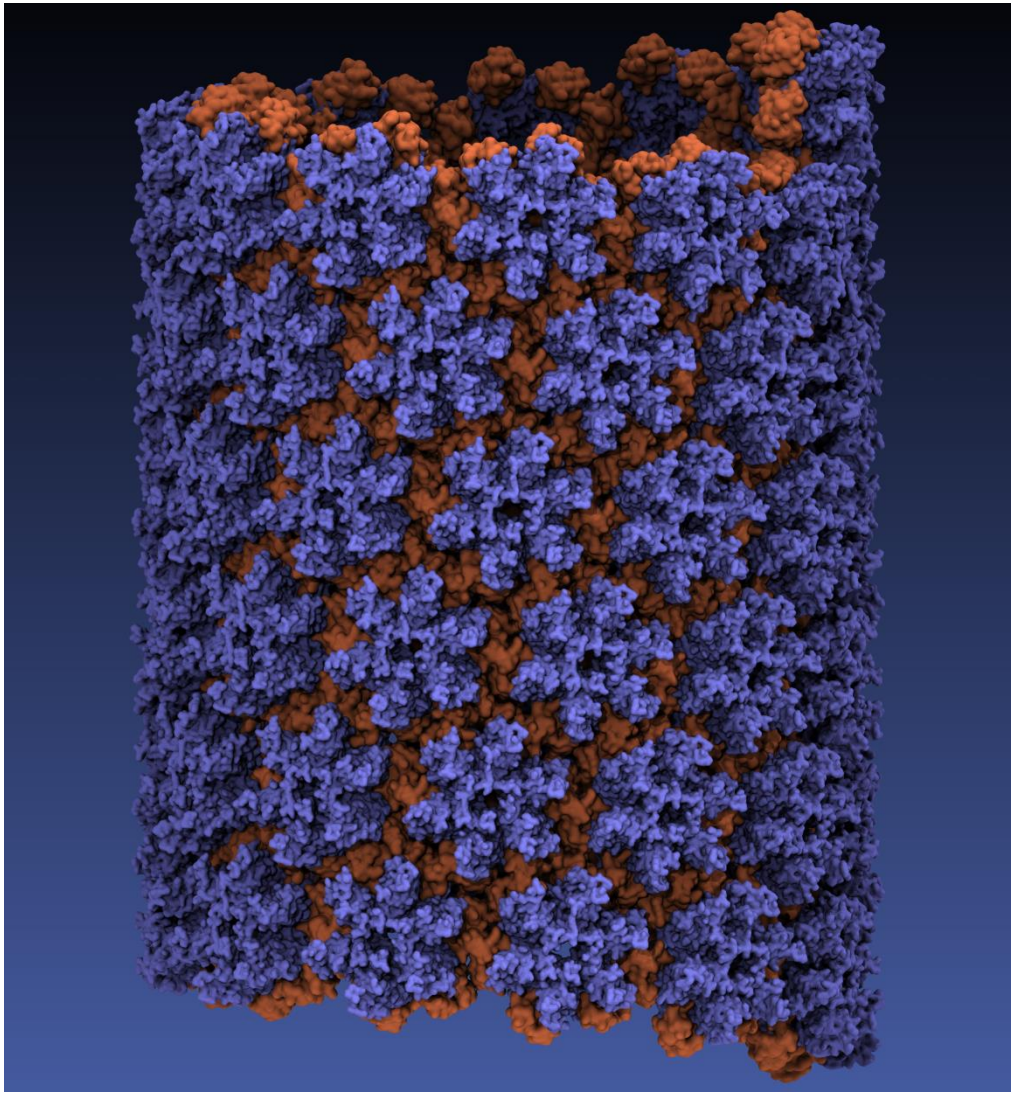
M. Krone, J. E. Stone, T. Ertl, K. Schulten. *EuroVis Short Papers*, pp. 67-71, 2012

Extensions and Analysis Uses for QuickSurf Triangle Mesh

- Curved PN triangles:
  - We have performed tests with post-processing the resulting triangle mesh and using curved PN triangles to generate smooth surfaces with a larger grid spacing, for increased performance
  - Initial results demonstrate some potential, but there can be pathological cases where MC generates long skinny triangles, causing unsightly surface creases
- Analysis uses (beyond visualization):
  - Minor modifications to the density map algorithm allow rapid computation of solvent accessible surface area by summing the areas in the resulting triangle mesh
  - Modifications to the density map algorithm will allow it to be used for MDFF (molecular dynamics flexible fitting)
  - Surface triangle mesh can be used as the input for computing the electrostatic potential field for mesh-based algorithms

Challenge: Support Interactive QuickSurf for Large Structures on Mid-Range GPUs

- Structures such as HIV initially needed large (6GB) GPU memory to generate fully-detailed surface renderings
- Goals and approach:
  - **Avoid slow CPU-fallback!**
  - Incrementally change algorithm phases to use more compact data types, while maintaining performance
  - Specialize code for different performance/memory capacity cases



Improving QuickSurf Memory Efficiency

- Both host and GPU memory capacity limitations are a significant concern when rendering surfaces for virus structures such as HIV or for large cellular models which can contain hundreds of millions of particles
- The original QuickSurf implementation used single-precision floating point for output vertex arrays and textures
- Judicious use of reduced-precision numerical representations, cut the overall memory footprint of the entire QuickSurf algorithm to half of the original
  - Data type changes made throughout the entire chain from density map computation through all stages of Marching Cubes

Supporting Multiple Data Types for QuickSurf Density Maps and Marching Cubes Vertex Arrays

- The major algorithm components of QuickSurf are now used for many other purposes:
  - Gaussian density map algorithm now used for MDFF Cryo EM density map fitting methods in addition to QuickSurf
  - Marching Cubes routines also used for Quantum Chemistry visualizations of molecular orbitals
- Rather than simply changing QuickSurf to use a particular internal numerical representation, it is desirable to instead use **CUDA C++ templates** to make type-generic versions of the key objects, kernels, and output vertex arrays
- Accuracy-sensitive algorithms use high-precision data types, performance and memory capacity sensitive cases use quantized or reduced precision approaches

Minimizing the Impact of Generality on QuickSurf Code Complexity

- A critical factor in the simplicity of supporting multiple QuickSurf data types arises from the so-called ***"gather"*** oriented algorithm we employ
  - Internally, all in-register arithmetic is single-precision
  - Data conversions to/from compressed or reduced precision data types are performed on-the-fly as needed
- Small **inlined** type conversion routines are defined for each of the cases we want to support
- Key QuickSurf kernels are genericized using C++ template syntax, and the compiler "connects the dots" to automatically generate type-specific kernels as needed

Example Templated Density Map Kernel

```
template<class DENSITY, class VOLTEX>
__global__ static void
gaussdensity_fast_tex_norm(int natoms,
                           const float4 * RESTRICT sorted_xyzr,
                           const float4 * RESTRICT sorted_color,
                           int3 numvoxels,
                           int3 acncells,
                          float acgridspacing,
                          float invacgridspacing,
                          const uint2 * RESTRICT cellStartEnd,
                          float gridspacing, unsigned int z,
                          DENSITY * RESTRICT densitygrid,
                          VOLTEX * RESTRICT voltexmap,
                         float invisovalue) {
```

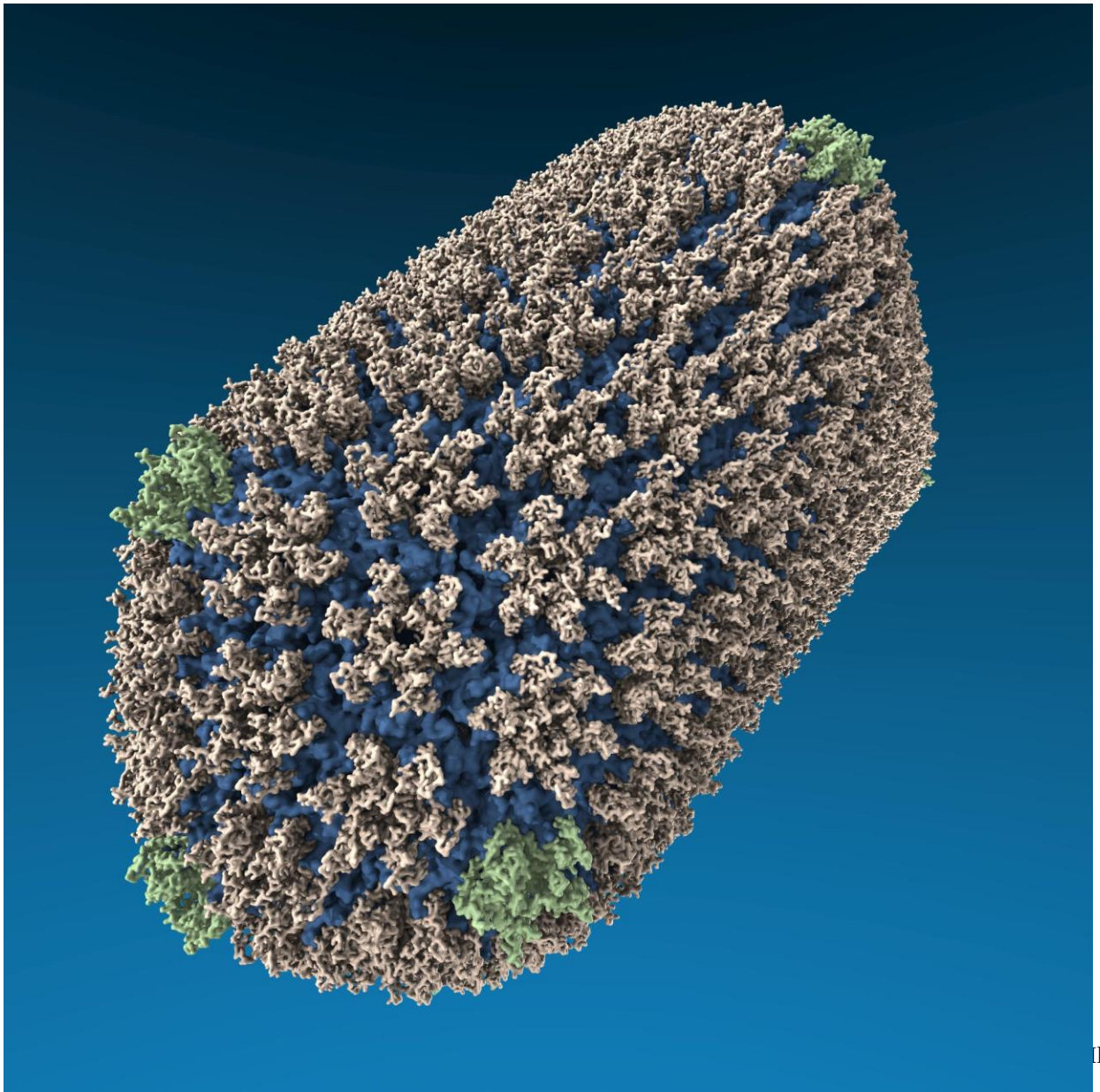
Example Templated Density Map Kernel

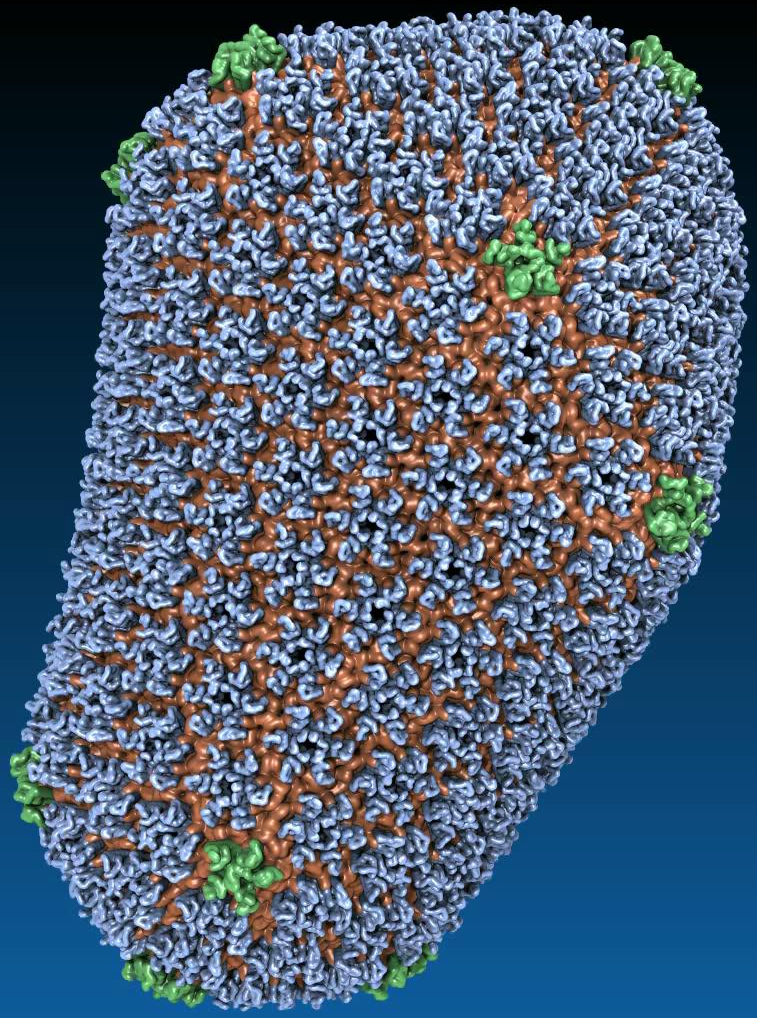
```
template<class DENSITY, class VOLTEX>
__global__ static void
gaussdensity_fast_tex_norm( …  ) {

  … Triple-nested and unrolled inner loops here …

  DENSITY densityout;
  VOLTEX texout;
  convert_density(densityout, densityval1);
  densitygrid[outaddr         ] = densityout;
  convert_color(texout, densitycol1);
  voltexmap[outaddr         ] = texout;
```

Net Result of QuickSurf Memory Efficiency Optimizations

- **Halved** overall GPU memory use
- Achieved **1.5x to 2x performance gain**:
  - The "gather" density map algorithm keeps type conversion operations out of the innermost loop
  - Density map global memory writes reduced to half
  - Multiple stages of Marching Cubes operate on smaller input and output data types
  - Same code path supports multiple precisions
- Users now get full GPU-accelerated QuickSurf in many cases that previously triggered CPU-fallback, all platforms (laptop/desk/super) benefit!

High Resolution HIV Surface



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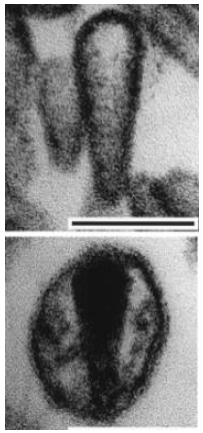
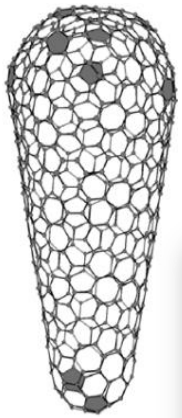
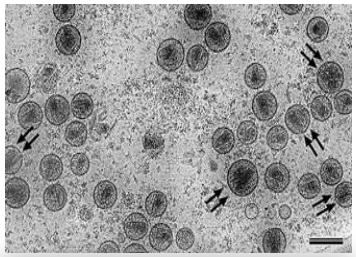
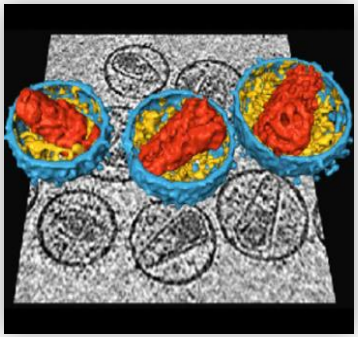
Structural Route to the all-atom HIV-1 Capsid

1st TEM (1999) 1st tomography (2003)









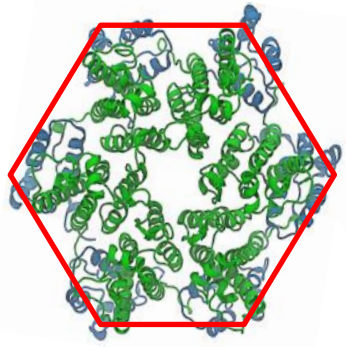
Ganser et al. *Science*, 1999

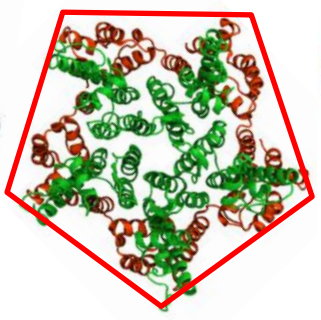
Briggs et al. *EMBO J*, 2003

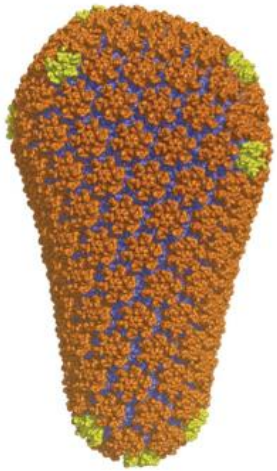
Briggs et al. *Structure*, 2006

cryo-ET (2006)

Crystal structures of separated hexamer and pentamer



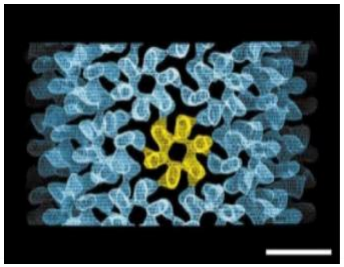




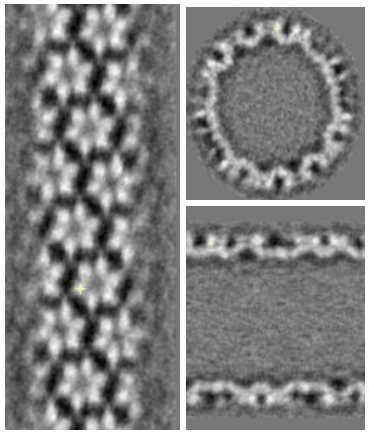
Pornillos et al. , *Cell* 2009, *Nature* 2011

High res. EM of hexameric tubule, tomography of capsid, **all-atom model of capsid by MDFF w/ NAMD & VMD, NSF/NCSA Blue Waters computer at Illinois**

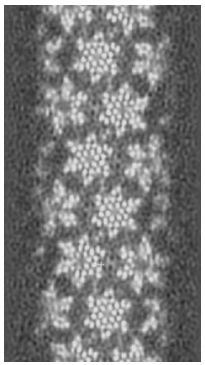
hexameric tubule

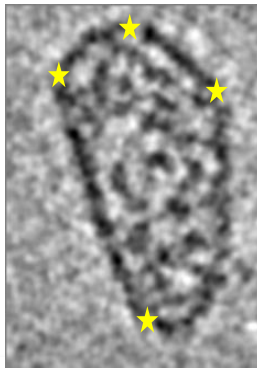


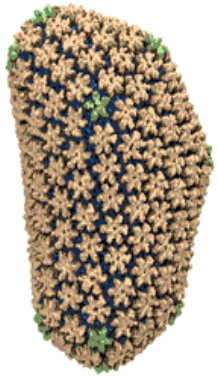
Li et al., *Nature*, 2000



Byeon et al., *Cell* 2009





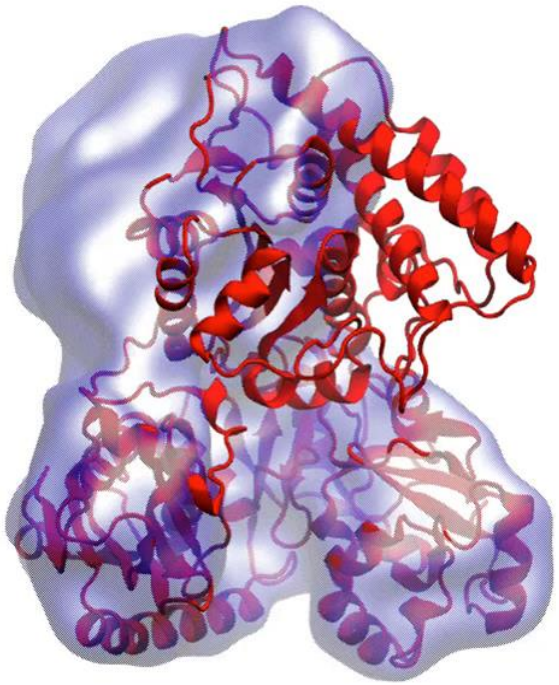


Zhao et al. , *Nature* 497: 643-646 (2013)

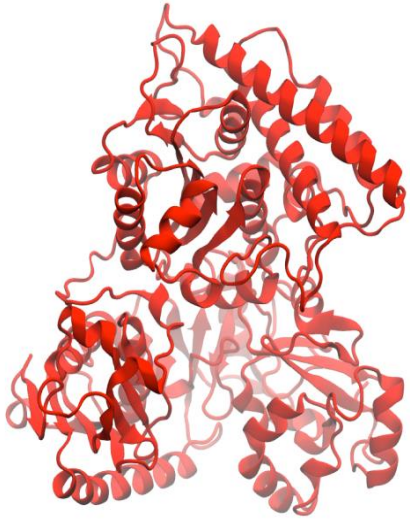
Molecular Dynamics Flexible Fitting (MDFF)

X-ray crystallography → MDFF ← Electron microscopy

APS at Argonne

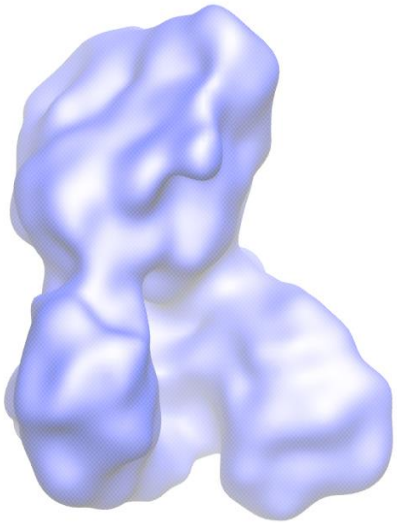


FEI microscope



Acetyl - CoA Synthase

ORNL Titan



Flexible fitting of atomic structures into electron microscopy maps using molecular dynamics.

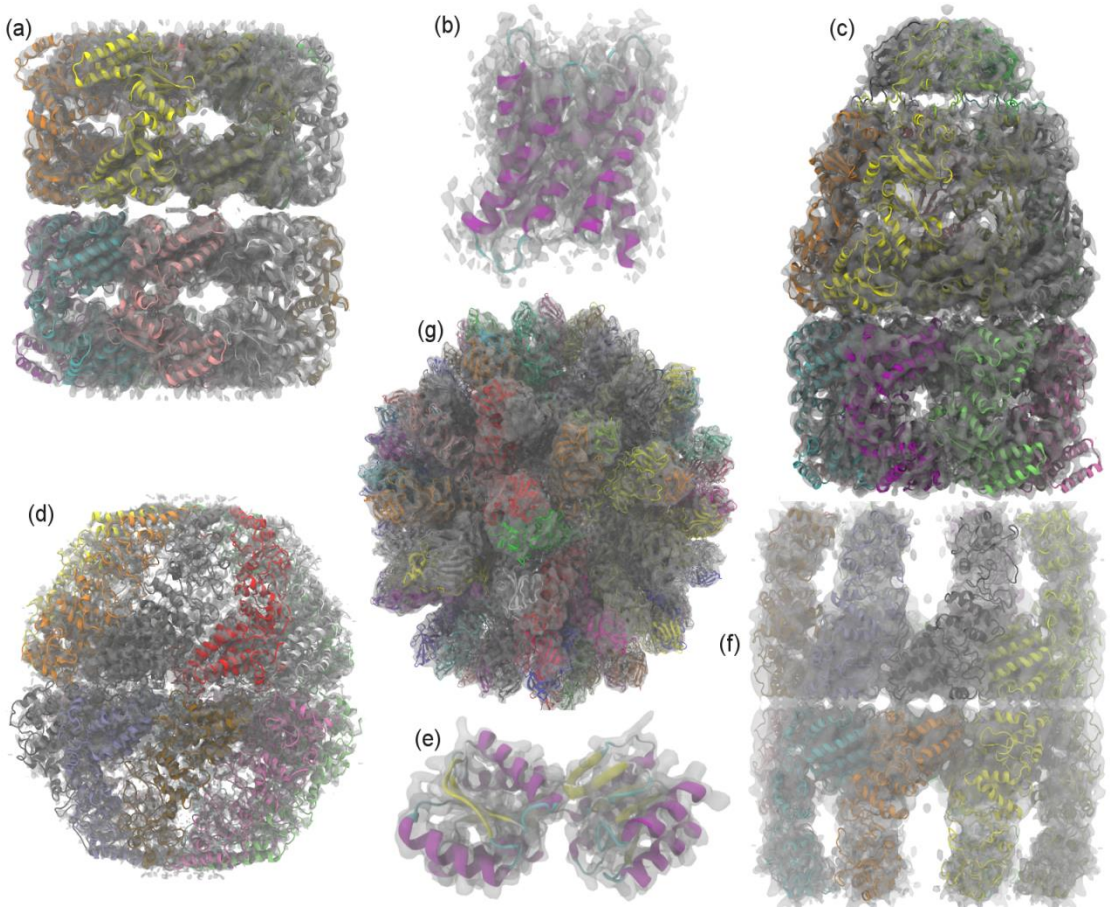
NIH BTRC for Macromolecular Modeling and Bioinformatics
http://www.ks.uiuc.edu/

Beckman Institute,
U. Illinois at Urbana-Champaign

L. Trabuco, E. Villa, K. Mitra, J. Frank, and K. Schulten. Structure, 16:673-683, 2008.

Evaluating Quality-of-Fit for Structures Solved by Hybrid Fitting Methods

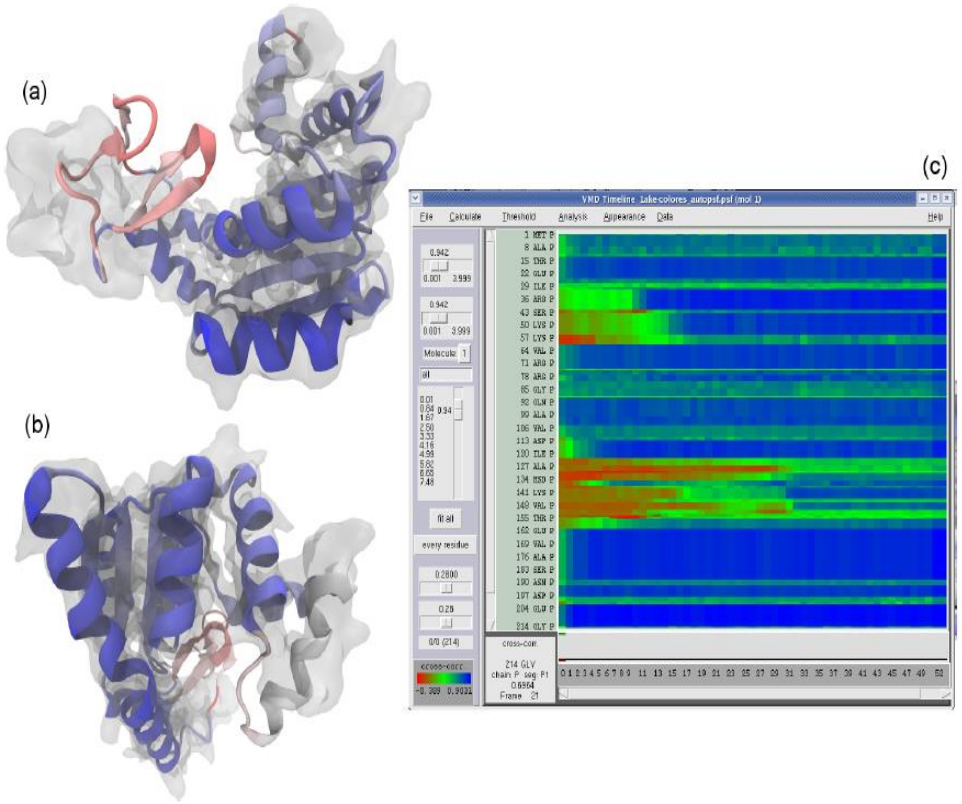
Compute Pearson correlation to evaluate the fit of a reference cryo-EM density map with a simulated density map produced from an all-atom structure.



GPUs Can Reduce Trajectory Analysis Runtimes from Hours to Minutes

GPUs enable laptops and desktop workstations to handle tasks that would have previously required a cluster, or a very long wait…

GPU-accelerated petascale supercomputers enable analyses were previously impossible, allowing detailed study of very large structures such as viruses



GPU-accelerated MDFF Cross Correlation Timeline

Regions with poor fit **Regions with good fit**

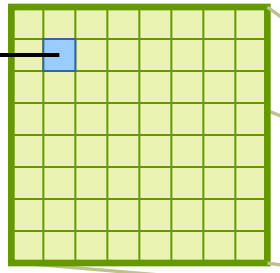
Single-Pass MDFF GPU Cross-Correlation

3-D density map decomposes into 3-D grid of 8x8x8 tiles containing CC partial sums and local CC values

Spatial CC map and overall CC value computed in a single pass

Small 8x8x2 CUDA thread blocks afford large per-thread register count, shared memory

Each thread computes 4 z-axis density map lattice points and associated CC partial sums



Threads producing results that are used

Inactive threads, region of discarded output

Grid of thread blocks

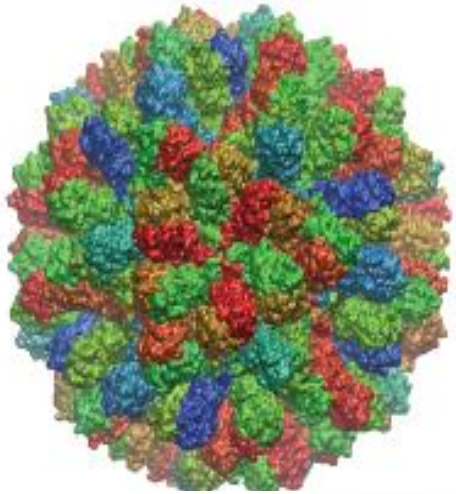
Padding optimizes global memory performance, guaranteeing coalesced global memory accesses

VMD GPU Cross Correlation Performance

	RHDV	Mm-cpn open	GroEL	Aquaporin
Resolution (Å)	**6.5**	**8**	**4**	**3**
Atoms	**702K**	**61K**	**54K**	**1.6K**
VMD-CUDA Quadro K6000	**0.458s 34.6x**	**0.06s 25.7x**	**0.034s 36.8x**	**0.007s 55.7x**
VMD-CPU-SSE 32-threads, 2x Xeon E5-2687W	**0.779s 20.3x**	**0.085s 18.1x**	**0.159s 7.9x**	**0.033s 11.8x**
Chimera 1-thread Xeon E5-2687W	**15.86s 1.0x**	**1.54s 1.0x**	**1.25s 1.0x**	**0.39s 1.0x**

GPU-accelerated analysis and visualization of large structures solved by molecular dynamics flexible fitting. J. E. Stone, R. McGreevy, B. Isralewitz, and K. Schulten. Faraday Discussion 169, 2014. (In press).

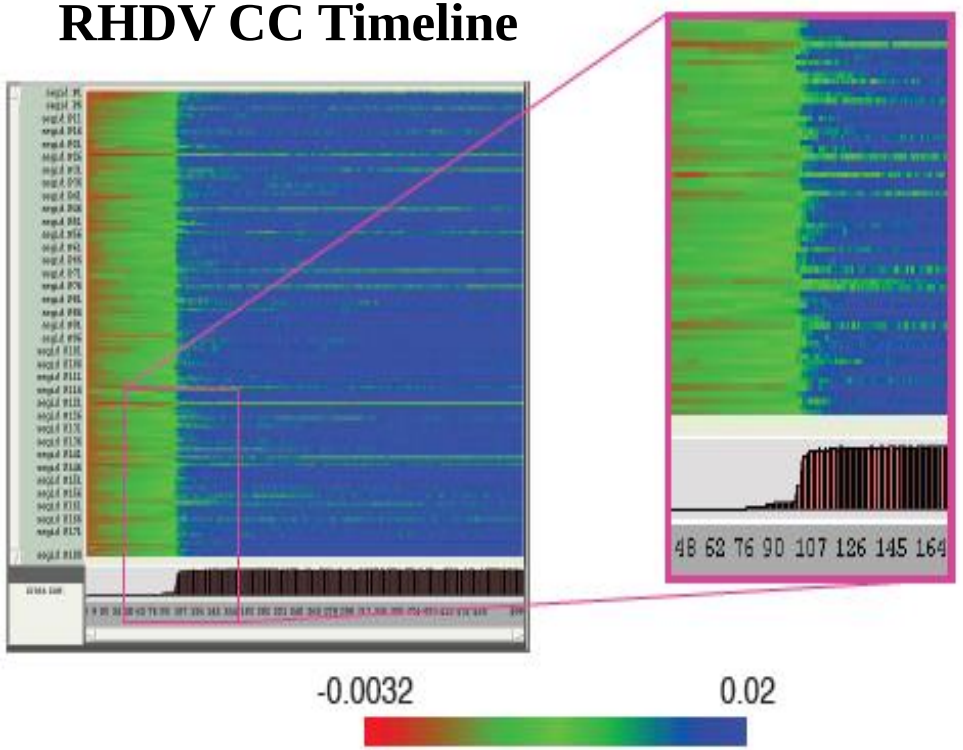
VMD RHDV Cross Correlation Timeline on Cray XK7



	RHDV
Atoms	**702K**
Component Selections	**720**
Single-node XK7 (projected)	**336 hours (14 days)**
128-node XK7	**3.2 hours 105x speedup**

Calculation would take **5 years** using conventional non-GPU software on a workstation!!

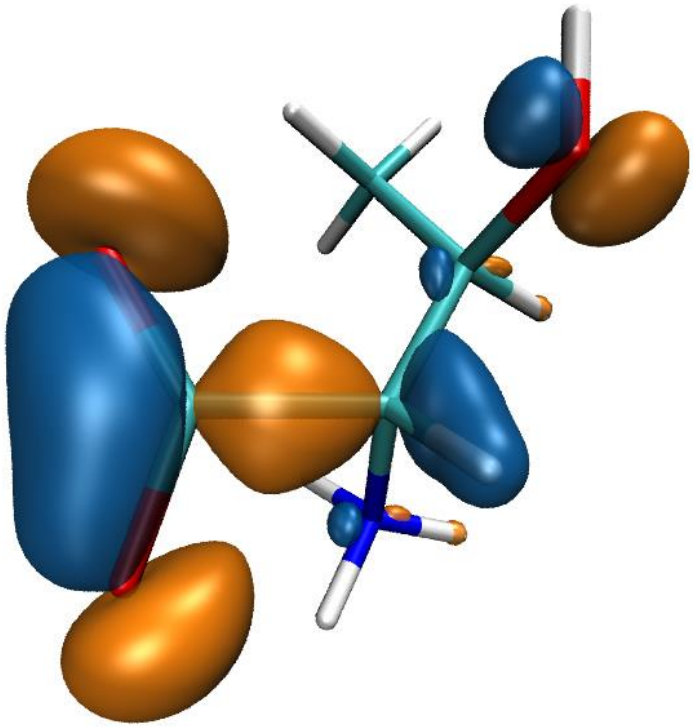
RHDV CC Timeline




1867

Molecular Orbitals

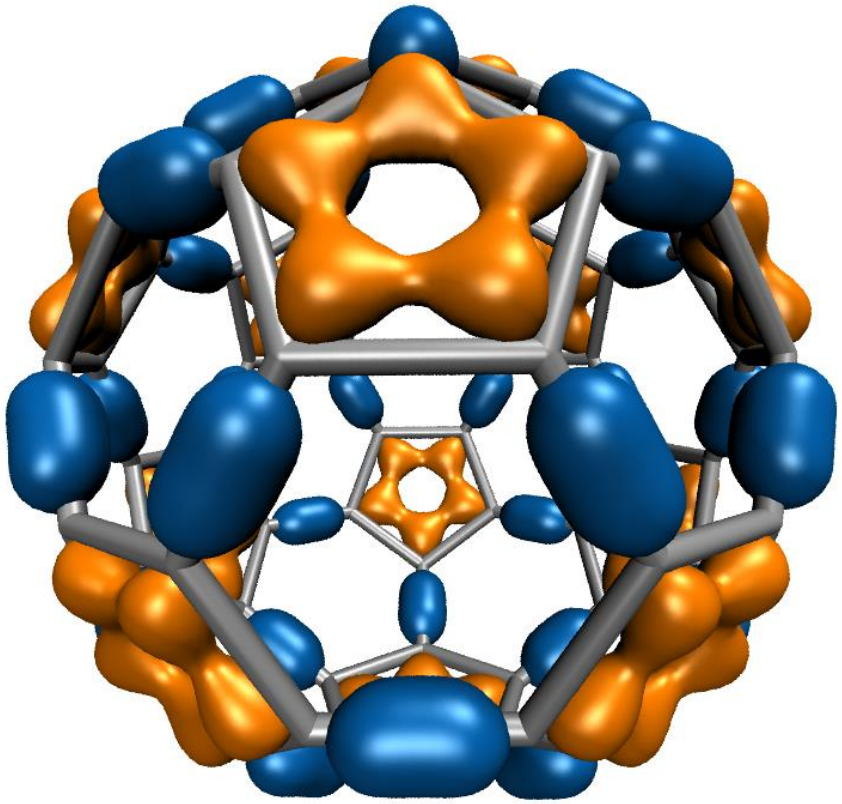
- Visualization of MOs aids in understanding the chemistry of molecular system
- MO spatial distribution is correlated with probability density for an electron(s)
- Algorithms for computing other molecular properties are similar, and can share code



High Performance Computation and Interactive Display of Molecular Orbitals on GPUs and Multi-core CPUs.
J. Stone, J. Saam, D. Hardy, K. Vandivort, W. Hwu, K. Schulten, *2nd Workshop on General-Purpose Computation on Graphics Pricessing Units (GPGPU-2), ACM International Conference Proceeding Series*, volume 383, pp. 9-18, 2009.

Computing Molecular Orbitals

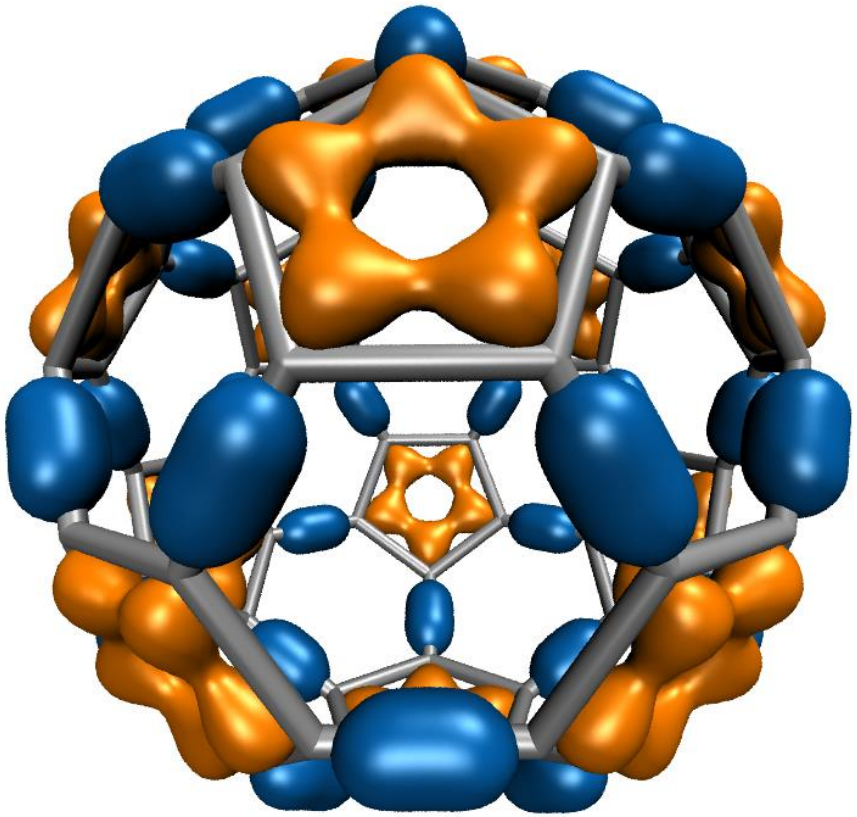
- Calculation of high resolution MO grids can require tens to hundreds of seconds in existing tools
- Existing tools cache MO grids as much as possible to avoid recomputation:
  - Doesn't eliminate the wait for initial calculation, hampers interactivity
  - Cached grids consume 100x-1000x more memory than MO coefficients



C_{60}

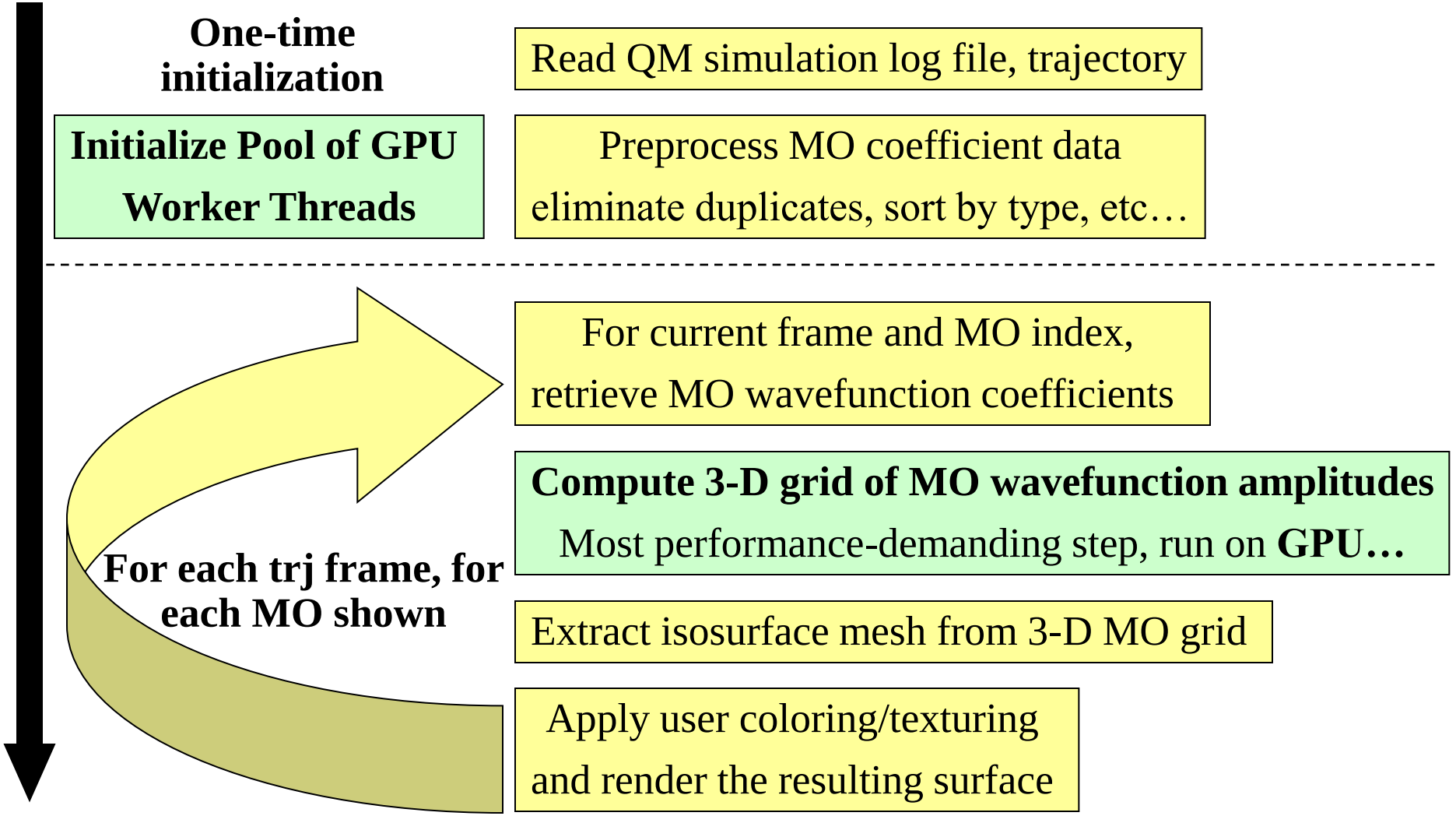
Animating Molecular Orbitals

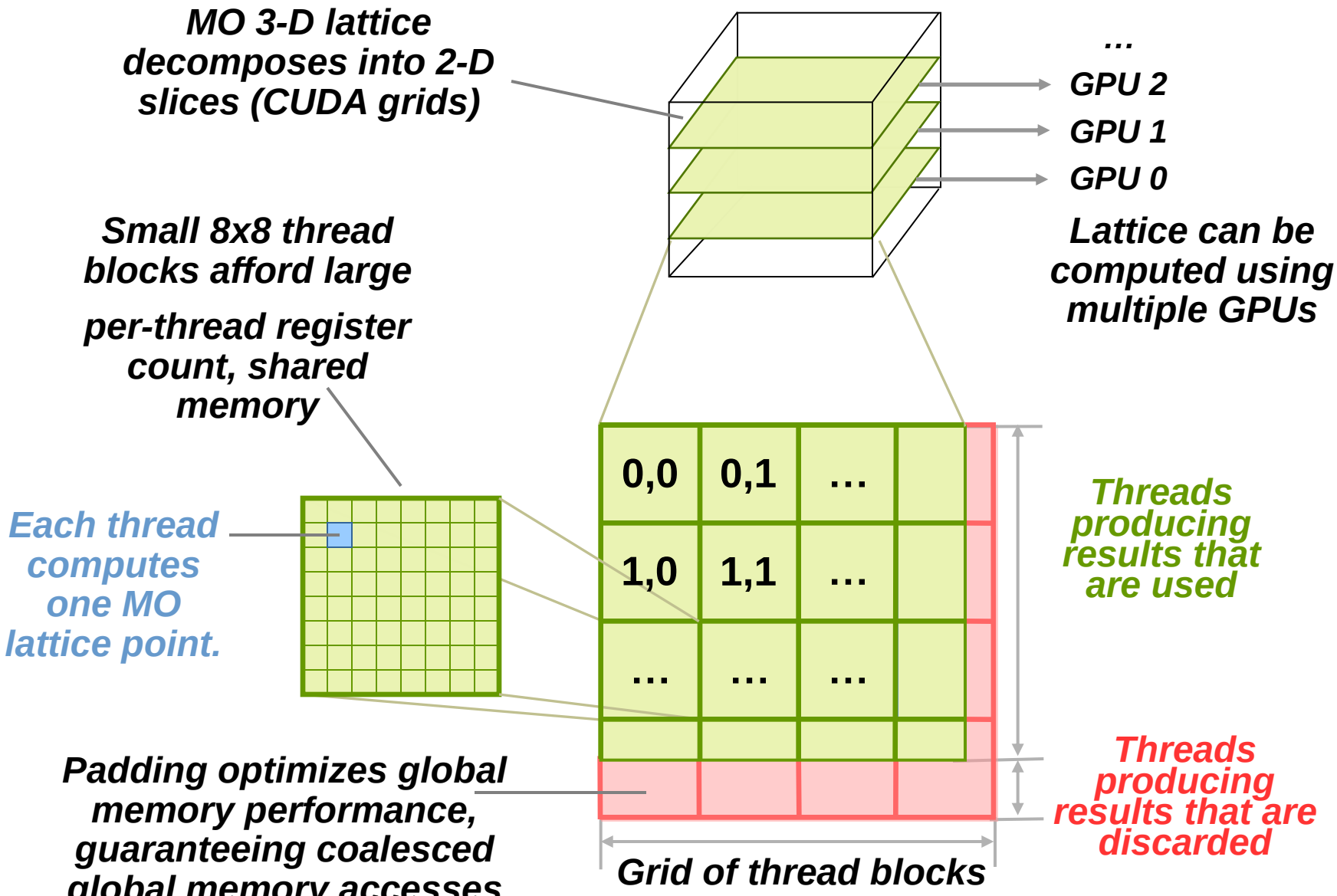
- Animation of (classical mechanics) molecular dynamics trajectories provides insight into simulation results
- To do the same for QM or QM/MM simulations one must compute MOs at ~**10 FPS** or more
- **>100x** speedup (GPU) over existing tools now makes this possible!



C_{60}

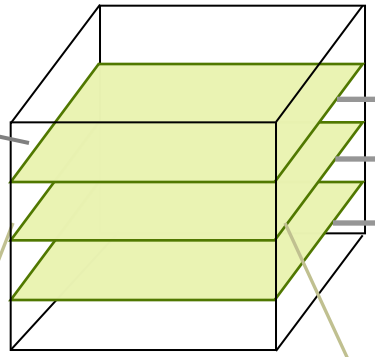
Molecular Orbital Computation and Display Process




MO 3-D lattice decomposes into 2-D slices (CUDA grids)
...
GPU 2
GPU 1
GPU 0
Lattice can be computed using multiple GPUs
Small 8x8 thread blocks afford large per-thread register count, shared memory
Each thread computes one MO lattice point.
0,0
0,1
…
1,0
1,1
…
…
…
…
Threads producing results that are used
Threads producing results that are discarded
Padding optimizes global memory performance, guaranteeing coalesced global memory accesses
Grid of thread blocks

MO GPU Parallel Decomposition

MO 3-D lattice decomposes into 2-D slices (CUDA grids)



...

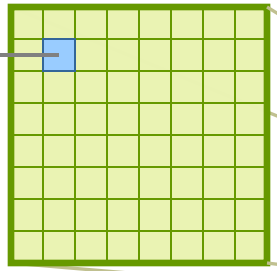
GPU 2

GPU 1

GPU 0

Lattice can be computed using multiple GPUs

Small 8x8 thread blocks afford large per-thread register count, shared memory



Each thread computes one MO lattice point.

Threads producing results that are used

Threads producing results that are discarded

Grid of thread blocks

Padding optimizes global memory performance, guaranteeing coalesced global memory accesses

MO GPU Kernel Snippet:
Contracted GTO Loop, Use of Constant Memory

```
[… outer loop over atoms …]
  float dist2 = xdist2 + ydist2 + zdist2;
  // Loop over the shells belonging to this atom (or basis function)
  for (shell=0; shell < maxshell; shell++) {
    float contracted_gto = 0.0f;
    // Loop over the Gaussian primitives of this contracted basis function to build the atomic
     orbital
    int maxprim = const_num_prim_per_shell[shell_counter];
    int shelltype = const_shell_types[shell_counter];
    for (prim=0; prim < maxprim;  prim++) {
      float exponent       = const_basis_array[prim_counter    ];
      float contract_coeff = const_basis_array[prim_counter + 1];
      contracted_gto += contract_coeff * __expf(-exponent*dist2);
      prim_counter += 2;
    }
[… continue on to angular momenta loop …]
```

Constant memory: nearly register-speed when array elements accessed in unison by all threads….

MO GPU Kernel Snippet:
Unrolled Angular Momenta Loop

```
/* multiply with the appropriate wavefunction coefficient */
float tmpshell=0;
switch (shelltype) {
  case S_SHELL:
    value += const_wave_f[ifunc++] * contracted_gto;
    break;
[… P_SHELL case …]
  case D_SHELL:
    tmpshell += const_wave_f[ifunc++] * xdist2;
    tmpshell += const_wave_f[ifunc++] * xdist * ydist;
    tmpshell += const_wave_f[ifunc++] * ydist2;
    tmpshell += const_wave_f[ifunc++] * xdist * zdist;
    tmpshell += const_wave_f[ifunc++] * ydist * zdist;
    tmpshell += const_wave_f[ifunc++] * zdist2;
    value += tmpshell * contracted_gto;
    break;
[... Other cases: F_SHELL, G_SHELL, etc …]
} // end switch
```

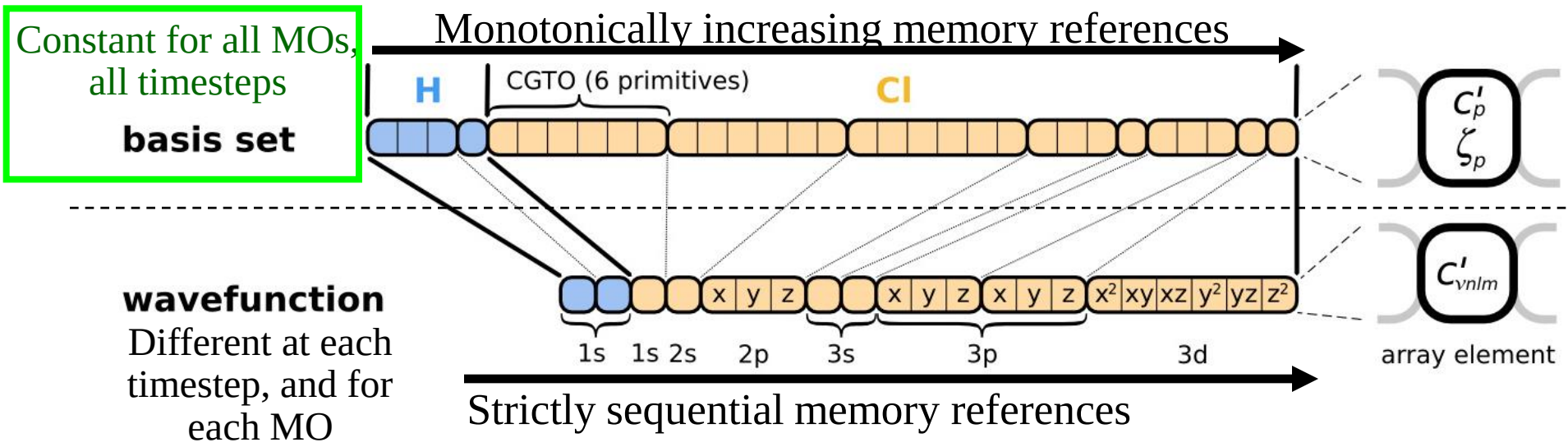
Loop unrolling:

- Saves registers (important for GPUs!)
- Reduces loop control overhead
- Increases arithmetic intensity

Preprocessing of Atoms, Basis Set, and Wavefunction Coefficients

- Must make effective use of high bandwidth, low-latency GPU on-chip shared memory, or L1 cache:
  - Overall storage requirement reduced by eliminating duplicate basis set coefficients
  - Sorting atoms by element type allows re-use of basis set coefficients for subsequent atoms of identical type
- Padding, alignment of arrays guarantees coalesced GPU global memory accesses

GPU Traversal of Atom Type, Basis Set, Shell Type, and Wavefunction Coefficients

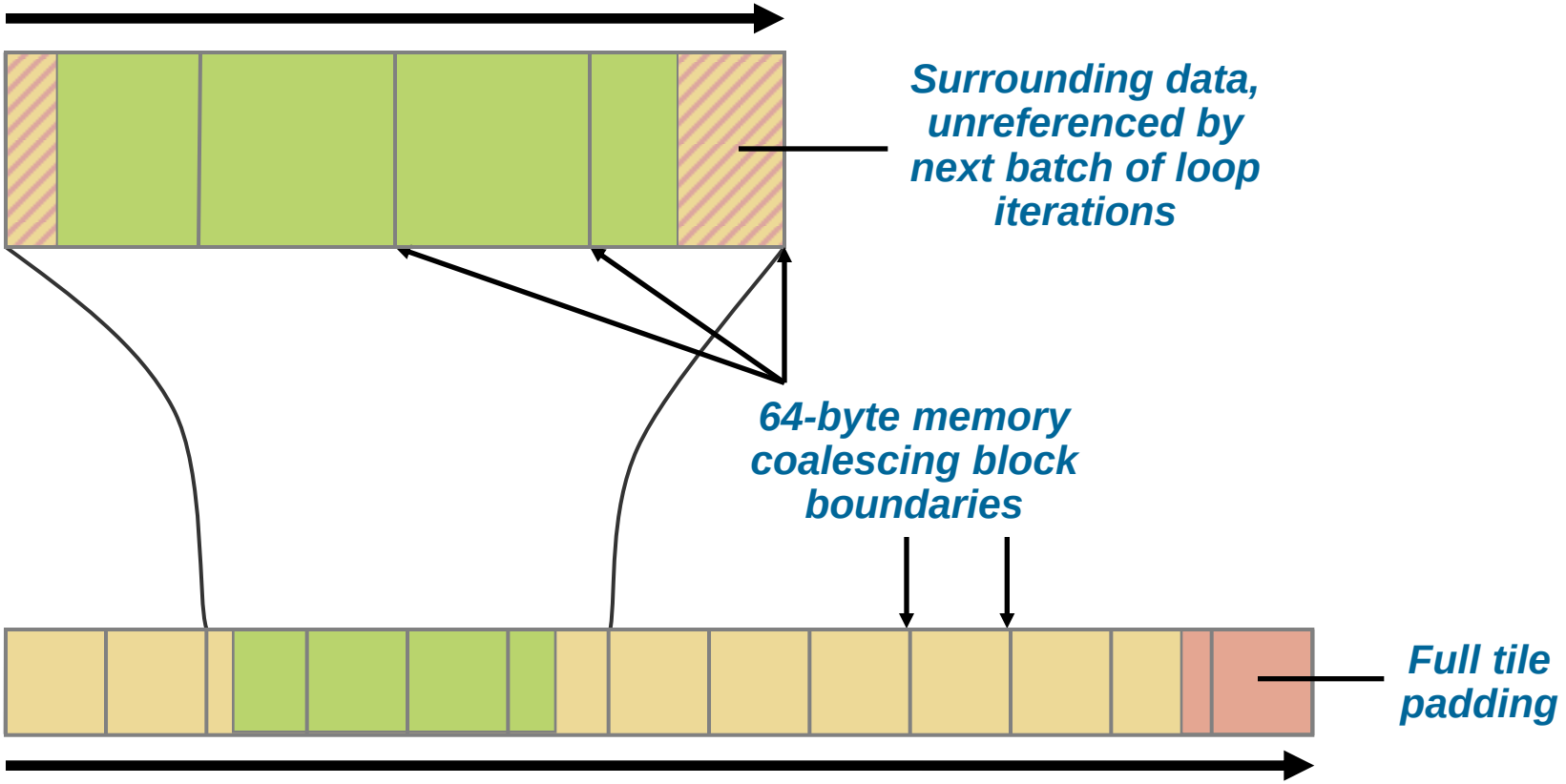


- Loop iterations always access same or consecutive array elements for all threads in a thread block:
  - Yields good constant memory and L1 cache performance
  - Increases shared memory tile reuse

Use of GPU On-chip Memory

- If total data less than 64 kB, use only const mem:
  - Broadcasts data to all threads, no global memory accesses!
- For large data, shared memory used as a program-managed cache, coefficients loaded on-demand:
  - Tiles sized large enough to service entire inner loop runs, broadcast to all 64 threads in a block
  - Complications: nested loops, multiple arrays, varying length
  - Key to performance is to locate tile loading checks outside of the two performance-critical inner loops
  - Only 27% slower than hardware caching provided by constant memory (on GT200)
- Fermi/Kepler GPUs have larger on-chip shared memory, L1/L2 caches, greatly reducing control overhead

Array tile loaded in GPU shared memory. Tile size is a power-of-two, a multiple of coalescing size, and allows simple indexing in inner loops. Global memory array indices are merely offset to reference an MO coefficient within a tile loaded in fast on-chip shared memory.



MO coefficient array in GPU global memory. Tiles are referenced in consecutive order.

VMD MO GPU Kernel Snippet:
Loading Tiles Into Shared Memory On-Demand

```
[… outer loop over atoms …]
  if ((prim_counter + (maxprim<<1)) >= SHAREDSIZE) {
    prim_counter += sblock_prim_counter;
    sblock_prim_counter = prim_counter & MEMCOAMASK;
    s_basis_array[sidx      ] = basis_array[sblock_prim_counter + sidx      ];
    s_basis_array[sidx +  64] = basis_array[sblock_prim_counter + sidx +  64];
    s_basis_array[sidx + 128] = basis_array[sblock_prim_counter + sidx + 128];
    s_basis_array[sidx + 192] = basis_array[sblock_prim_counter + sidx + 192];
    prim_counter -= sblock_prim_counter;
    __syncthreads();
  }
  for (prim=0; prim < maxprim;  prim++) {
    float exponent       = s_basis_array[prim_counter    ];
    float contract_coeff = s_basis_array[prim_counter + 1];
    contracted_gto += contract_coeff * __expf(-exponent*dist2);
    prim_counter += 2;
  }
[… continue on to angular momenta loop …]
```

Shared memory tiles:

- Tiles are checked and loaded, if necessary, immediately prior to entering key arithmetic loops
- Adds additional control overhead to loops, even with optimized implementation

New GPUs Bring Opportunities for Higher Performance and Easier Programming

- NVIDIA's "Fermi" and "Kepler" GPUs bring:
  - **Greatly increased** peak single- and double-precision arithmetic rates
  - **Moderately** increased global memory bandwidth
  - Increased capacity on-chip memory partitioned into shared memory and an L1 cache for global memory
  - Concurrent kernel execution
  - Bidirectional asynchronous host-device I/O
  - ECC memory, faster atomic ops, many others…

VMD MO GPU Kernel Snippet:
Fermi/Kepler kernel based on L1 cache

```
[… outer loop over atoms …]
  // loop over the shells/basis funcs belonging to this atom
  for (shell=0; shell < maxshell; shell++) {
    float contracted_gto = 0.0f;
    int maxprim    = shellinfo[(shell_counter<<4)     ];
    int shell_type = shellinfo[(shell_counter<<4) + 1];
    for (prim=0; prim < maxprim; prim++) {
      float exponent       = basis_array[prim_counter     ];
      float contract_coeff = basis_array[prim_counter + 1];
      contracted_gto += contract_coeff * __expf(-
      exponent*dist2);
      prim_counter += 2;
  }
  […continue on to angular momenta loop …]
```

L1 cache:

- Simplifies code!
- Reduces control overhead
- Gracefully handles arbitrary-sized problems
- Matches performance of constant memory on Fermi

MO Kernel for One Grid Point (Naive C)

```
…
for (at=0; at<numatoms; at++) {
  int prim_counter = atom_basis[at];
  calc_distances_to_atom(&atompos[at], &xdist, &ydist, &zdist, &dist2, &xdiv);
  for (contracted_gto=0.0f, shell=0; shell < num_shells_per_atom[at]; shell++) {
    int shell_type = shell_symmetry[shell_counter];
    for (prim=0; prim < num_prim_per_shell[shell_counter];  prim++) {
      float exponent       = basis_array[prim_counter    ];
      float contract_coeff = basis_array[prim_counter + 1];
      contracted_gto += contract_coeff * expf(-exponent*dist2);
      prim_counter += 2;
    }
    for (tmpshell=0.0f, j=0, zdp=1.0f; j<=shell_type; j++, zdp*=zdist) {
      int imax = shell_type - j;
      for (i=0, ydp=1.0f, xdp=pow(xdist, imax); i<=imax; i++, ydp*=ydist, xdp*=xdiv)
        tmpshell += wave_f[ifunc++] * xdp * ydp * zdp;
    }
    value += tmpshell * contracted_gto;
    shell_counter++;
  }
} …..
```

Loop over atoms

Loop over shells

Loop over primitives: largest component of runtime, due to expf()

Loop over angular momenta

(unrolled in real code)

Use of GPU On-chip Memory

- If total data less than 64 kB, use only const mem:
  - Broadcasts data to all threads, no global memory accesses!
- For large data, shared memory used as a program-managed cache, coefficients loaded on-demand:
  - Tile data in shared mem is broadcast to 64 threads in a block
  - Nested loops traverse multiple coefficient arrays of varying length, complicates things significantly…
  - Key to performance is to locate tile loading checks outside of the two performance-critical inner loops
  - Tiles sized large enough to service entire inner loop runs
  - Only 27% slower than hardware caching provided by constant memory (GT200)

Performance Evaluation:

Molekel, MacMolPlt, and VMD

Sun Ultra 24: Intel Q6600, NVIDIA GTX 280

	C_{60}-A	C_{60}-B	Thr-A	Thr-B	Kr-A	Kr-B
Atoms	60	60	17	17	1	1
Basis funcs (unique)	**300 (5)**	**900 (15)**	**49 (16)**	**170 (59)**	**19 (19)**	**84 (84)**

Kernel	Cores GPUs	Speedup vs. Molekel on 1 CPU core					
Molekel	1*	1.0	1.0	1.0	1.0	1.0	1.0
MacMolPlt	4	2.4	2.6	2.1	2.4	4.3	4.5
VMD GCC-cephes	4	3.2	4.0	3.0	3.5	4.3	6.5
VMD ICC-SSE-cephes	4	16.8	17.2	13.9	12.6	17.3	21.5
VMD ICC-SSE-approx**	4	59.3	53.4	50.4	49.2	54.8	69.8
VMD CUDA-const-cache	1	552.3	533.5	355.9	421.3	193.1	571.6

VMD MO Performance Results for C_{60}

Sun Ultra 24: Intel Q6600, NVIDIA GTX 280

Kernel	Cores/GPUs	Runtime (s)	Speedup
CPU ICC-SSE	1	46.58	1.00
CPU ICC-SSE	4	11.74	3.97
CPU ICC-SSE-approx**	4	3.76	12.4
CUDA-tiled-shared	1	0.46	100.
CUDA-const-cache	1	0.37	126.
CUDA-const-cache-JIT*	**1**	**0.27**	**173. (JIT 40% faster)**

C_{60} basis set 6-31Gd. We used an unusually-high resolution MO grid for accurate timings. A more typical calculation has $1/8^{th}$ the grid points.

* Runtime-generated JIT kernel compiled using batch mode CUDA tools

**Reduced-accuracy approximation of expf(), cannot be used for zero-valued MO isosurfaces

VMD Single-GPU Molecular Orbital Performance Results for C_{60} on Fermi

Intel X5550 CPU, GeForce GTX 480 GPU

Kernel	Cores/GPUs	Runtime (s)	Speedup
Xeon 5550 ICC-SSE	1	30.64	1.0
Xeon 5550 ICC-SSE	8	4.13	7.4
CUDA shared mem	1	0.37	83
CUDA L1-cache (16KB)	**1**	**0.27**	**113**
CUDA const-cache	1	0.26	117
CUDA const-cache, zero-copy	1	0.25	122

Fermi GPUs have caches: match perf. of hand-coded shared memory kernels. Zero-copy memory transfers improve overlap of computation and host-GPU I/Os.

Preliminary Single-GPU Molecular Orbital Performance Results for C_{60} on Kepler

Intel X5550 CPU, GeForce GTX 680 GPU

Kernel	Cores/GPUs	Runtime (s)	Speedup
Xeon 5550 ICC-SSE	1	30.64	1.0
Xeon 5550 ICC-SSE	8	4.13	7.4
CUDA shared mem	1	0.264	116
CUDA L1-cache (16KB)	1	0.228	134
CUDA const-cache	1	0.104	292
CUDA const-cache, zero-copy	1	0.0938	326

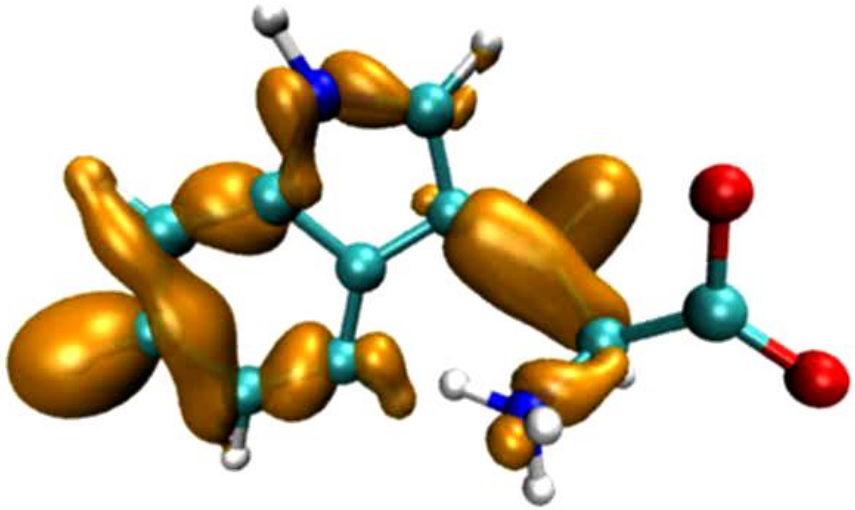
Kepler GK104 (GeForce 680) strongly prefers the constant cache kernels vs. the others.

VMD Orbital Dynamics Proof of Concept

One GPU can compute and animate this movie on-the-fly!

CUDA const-cache kernel, Sun Ultra 24, GeForce GTX 285

GPU MO grid calc.	**0.016 s**
CPU surface gen, volume gradient, and GPU rendering	**0.033 s**
Total runtime	**0.049 s**
Frame rate	**20 FPS**



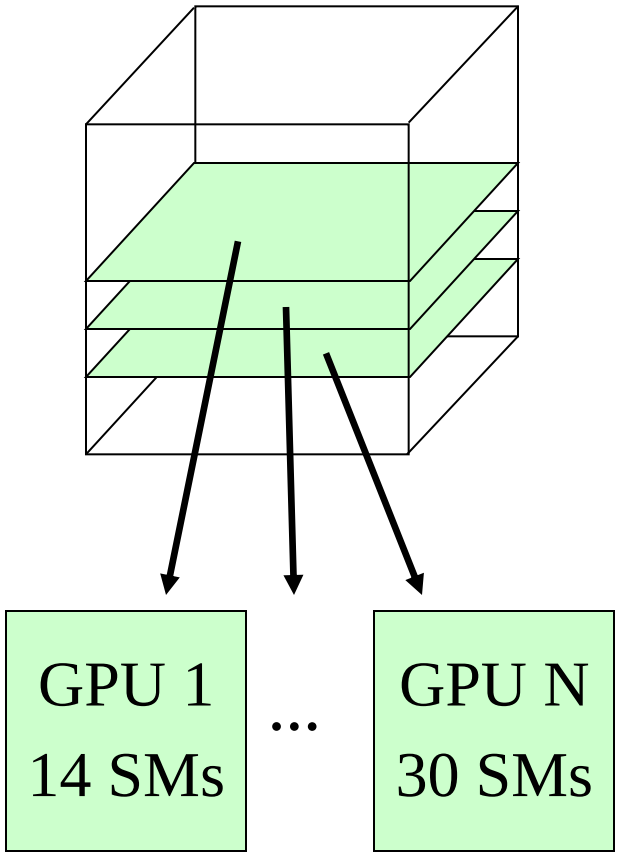
threonine

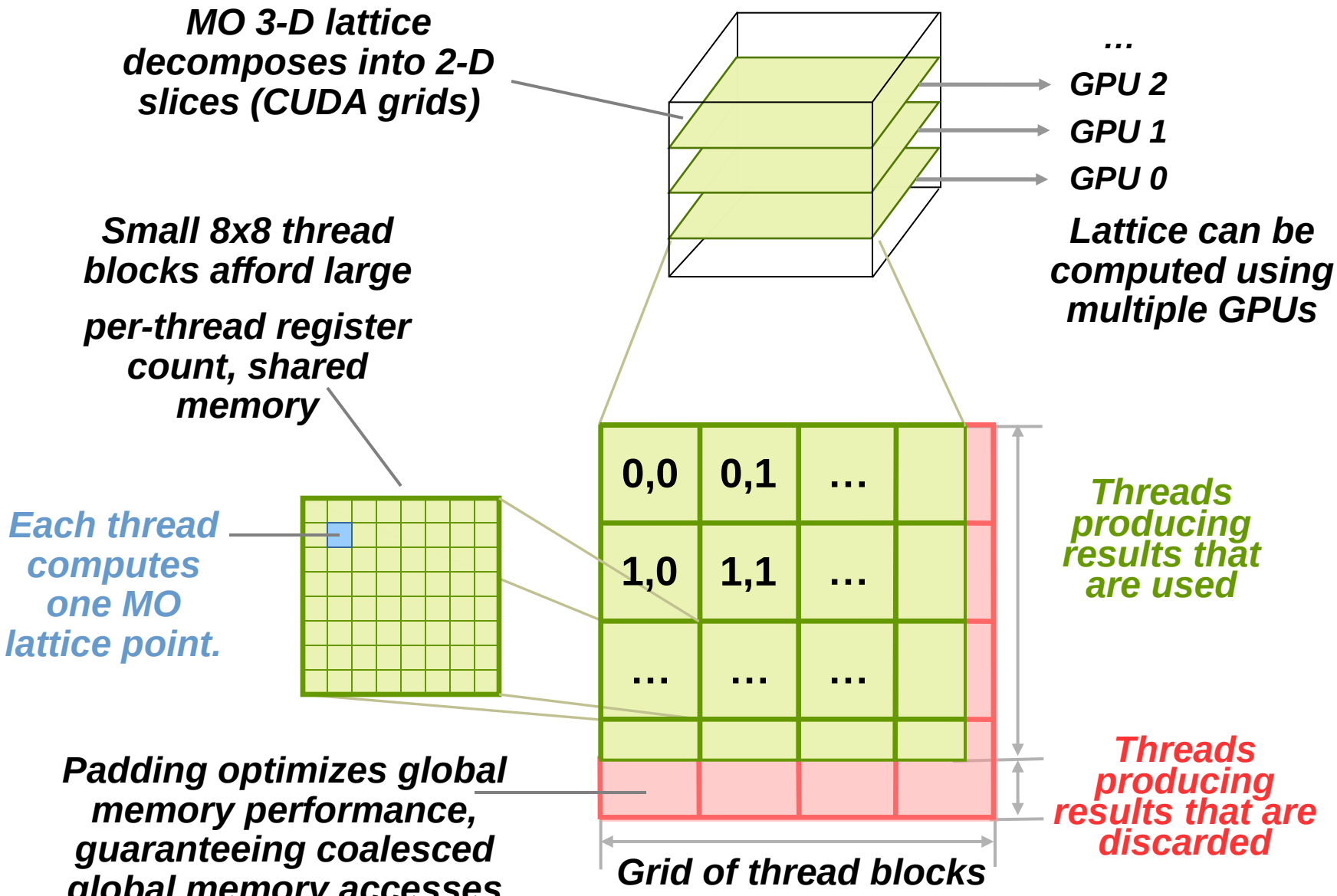
With GPU speedups over **100x**, previously insignificant CPU surface gen, gradient calc, and rendering are now **66%** of runtime.

Need GPU-accelerated surface gen next…

Multi-GPU Load Balance

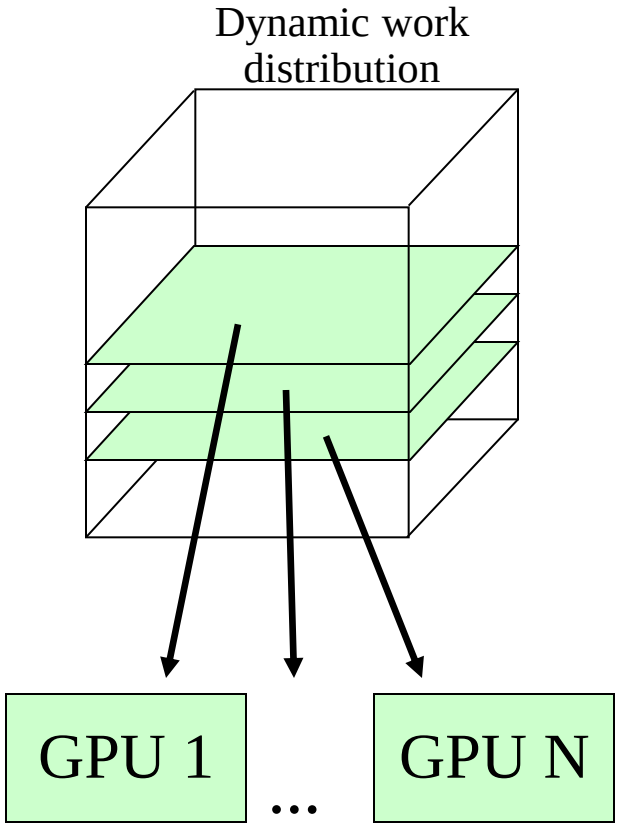
- Many early CUDA codes assumed all GPUs were identical
- Host machines may contain a diversity of GPUs of varying capability (discrete, IGP, etc)
- Different GPU on-chip and global memory capacities may need different problem “tile” sizes
- Static decomposition works poorly for non-uniform workload, or diverse GPUs




MO 3-D lattice decomposes into 2-D slices (CUDA grids)
...
GPU 2
GPU 1
GPU 0
Lattice can be computed using multiple GPUs
Small 8x8 thread blocks afford large per-thread register count, shared memory
Each thread computes one MO lattice point.
0,0
0,1
…
1,0
1,1
…
…
…
…
Threads producing results that are used
Threads producing results that are discarded
Padding optimizes global memory performance, guaranteeing coalesced global memory accesses
Grid of thread blocks

Multi-GPU Dynamic Work Distribution

```
// Each GPU worker thread loops over
// subset of work items…
while (!threadpool_next_tile(&parms,
    tilesize, &tile){
  // Process one work item…
  // Launch one CUDA kernel for each
  //   loop iteration taken…
  // Shared iterator automatically
  //   balances load on GPUs
}
```

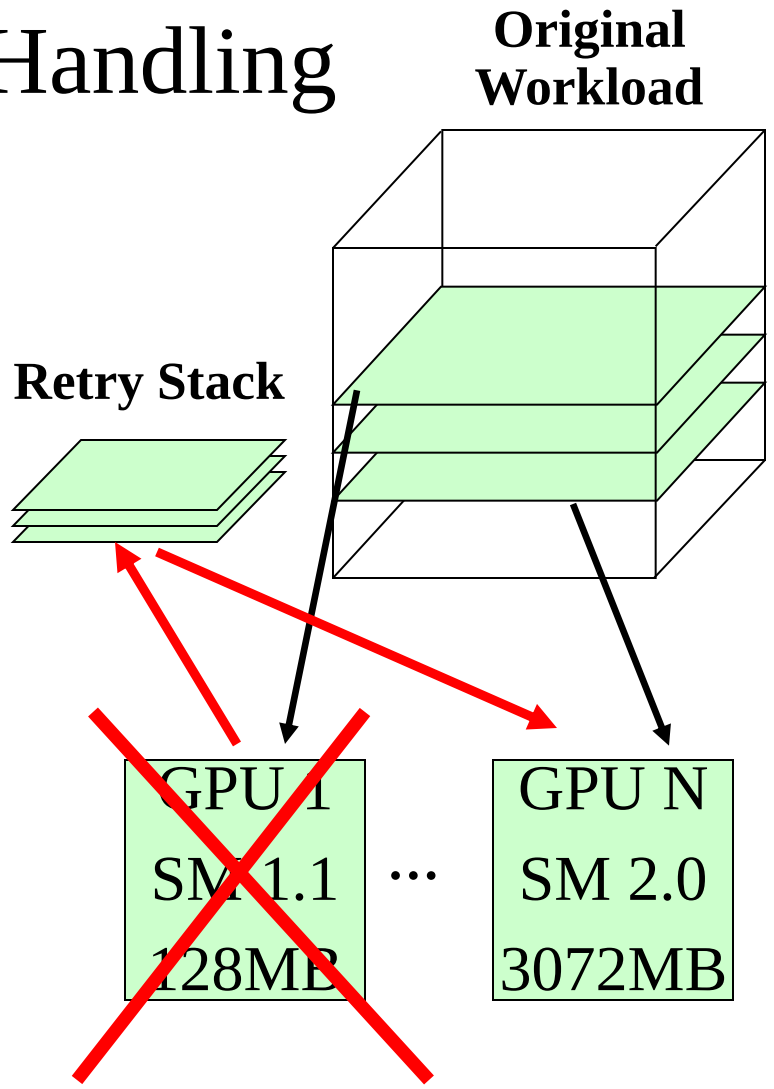


Example Multi-GPU Latencies
4 C2050 GPUs, Intel Xeon 5550

6.3us	CUDA empty kernel (immediate return)
9.0us	Sleeping barrier primitive (non-spinning barrier that uses POSIX condition variables to prevent idle CPU consumption while workers wait at the barrier)
14.8us	pool wake, host fctn exec, sleep cycle (no CUDA)
30.6us	pool wake, 1x(tile fetch, simple CUDA kernel launch), sleep
1817.0us	pool wake, 100x(tile fetch, simple CUDA kernel launch), sleep

Multi-GPU Runtime Error/Exception Handling

- Competition for resources from other applications can cause runtime failures, e.g. GPU out of memory half way through an algorithm
- Handle exceptions, e.g. convergence failure, NaN result, insufficient compute capability/features
- Handle and/or reschedule failed tiles of work



VMD Multi-GPU Molecular Orbital Performance Results for C_{60}

Kernel	Cores/GPUs	Runtime (s)	Speedup	Parallel Efficiency
CPU-ICC-SSE	1	46.580	1.00	100%
CPU-ICC-SSE	4	11.740	3.97	99%
CUDA-const-cache	1	0.417	112	100%
CUDA-const-cache	2	0.220	212	94%
CUDA-const-cache	3	0.151	308	92%
CUDA-const-cache	4	0.113	412	92%

Intel Q6600 CPU, 4x Tesla C1060 GPUs,

Uses persistent thread pool to avoid GPU init overhead, dynamic scheduler distributes work to GPUs

VMD Multi-GPU Molecular Orbital Performance Results for C_{60}

Intel X5550 CPU, 4x GeForce GTX 480 GPUs,

Kernel	Cores/GPUs	Runtime (s)	Speedup
Intel X5550-SSE	1	30.64	1.0
Intel X5550-SSE	8	4.13	7.4
GeForce GTX 480	1	0.255	120
GeForce GTX 480	2	0.136	225
GeForce GTX 480	3	0.098	312
GeForce GTX 480	4	0.081	378

Uses persistent thread pool to avoid GPU init overhead, dynamic scheduler distributes work to GPUs

Molecular Orbital Dynamic Scheduling Performance with Heterogeneous GPUs

Kernel	Cores/GPUs	Runtime (s)	Speedup
Intel X5550-SSE	1	30.64	1.0
Quadro 5800	1	0.384	79
Tesla C2050	1	0.325	94
GeForce GTX 480	1	0.255	120
GeForce GTX 480 + Tesla C2050 + Quadro 5800	3	0.114	268 (91% of ideal perf)

Dynamic load balancing enables mixture of GPU generations, SM counts, and clock rates to perform well.

MO Kernel Structure, Opportunity for JIT…

Data-driven, but representative loop trip counts in (…)

Loop over atoms (1 to ~200) {

Loop over electron shells for this atom type (1 to ~6) {

Loop over primitive functions for this shell type (1 to ~6) {

Unpredictable (at compile-time, since data-driven) but small loop trip counts result in significant loop overhead. **Dynamic kernel generation and JIT compilation can unroll entirely, resulting in 40% speed boost**

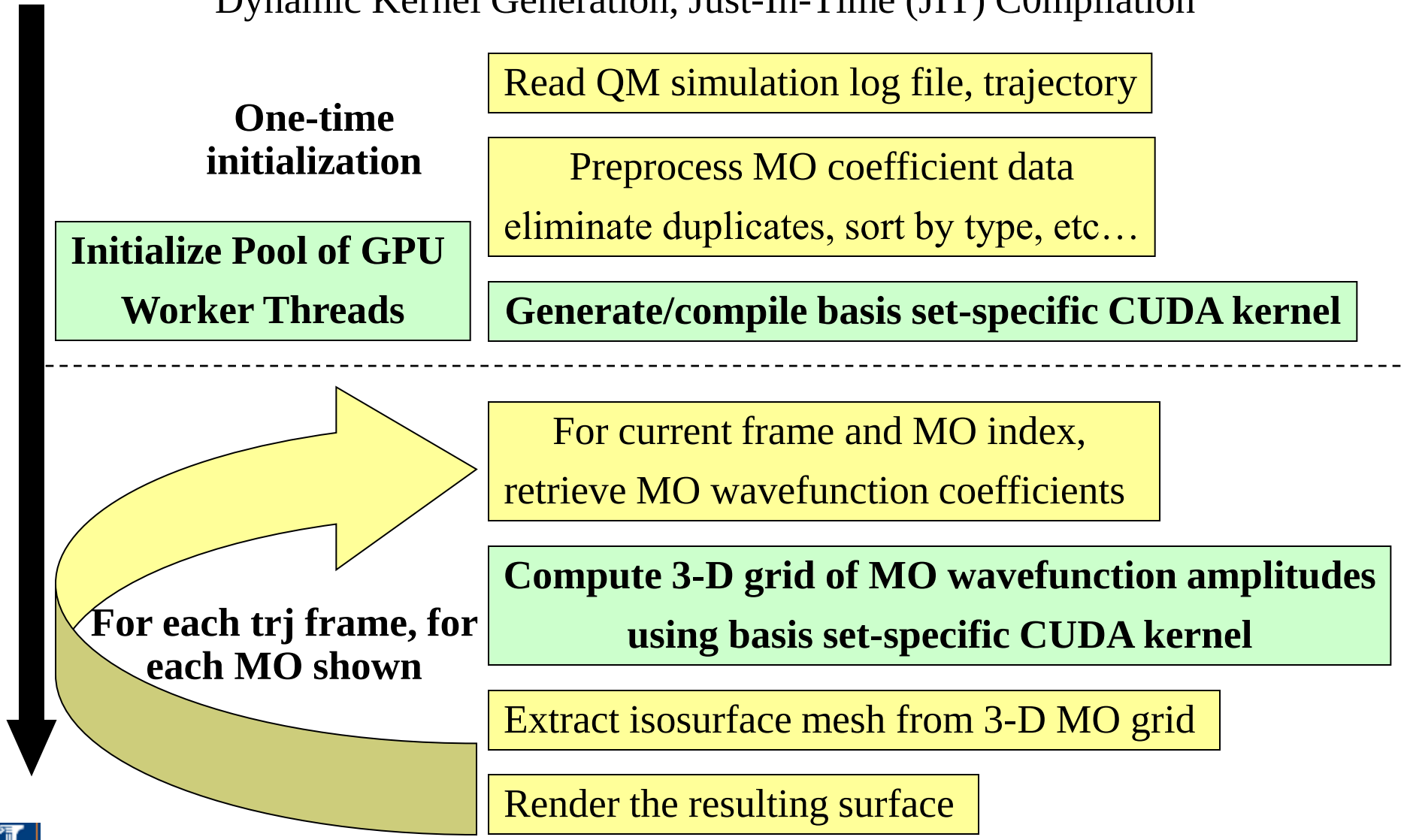
Loop over angular momenta for this shell type (1 to ~15) {}

}

}

Molecular Orbital Computation and Display Process

Dynamic Kernel Generation, Just-In-Time (JIT) C0mpilation



```
…..
// loop over the shells belonging to this atom (or basis function)
for (shell=0; shell < maxshell; shell++) {
  float contracted_gto = 0.0f;

  // Loop over the Gaussian primitives of this contracted
  // basis function to build the atomic orbital
  int maxprim = const_num_prim_per_shell[shell_counter];
  int shell_type = const_shell_symmetry[shell_counter];
  for (prim=0; prim < maxprim;  prim++) {
    float exponent       = const_basis_array[prim_counter    ];
    float contract_coeff = const_basis_array[prim_counter + 1];
    contracted_gto += contract_coeff * exp2f(-exponent*dist2);
    prim_counter += 2;
  }

  /* multiply with the appropriate wavefunction coefficient */
  float tmpshell=0;
  switch (shell_type) {
    case S_SHELL:
      value += const_wave_f[ifunc++] * contracted_gto;
      break;
[…..]
    case D_SHELL:
      tmpshell += const_wave_f[ifunc++] * xdist2;
      tmpshell += const_wave_f[ifunc++] * ydist2;
      tmpshell += const_wave_f[ifunc++] * zdist2;
      tmpshell += const_wave_f[ifunc++] * xdist * ydist;
      tmpshell += const_wave_f[ifunc++] * xdist * zdist;
      tmpshell += const_wave_f[ifunc++] * ydist * zdist;
      value += tmpshell * contracted_gto;
      break;
```

General loop-based CUDA kernel

Dynamically-generated CUDA kernel (JIT)

```
…..
contracted_gto = 1.832937 * expf(-7.868272*dist2);
contracted_gto += 1.405380 * expf(-1.881289*dist2);
contracted_gto += 0.701383 * expf(-0.544249*dist2);
// P_SHELL
tmpshell = const_wave_f[ifunc++] * xdist;
tmpshell += const_wave_f[ifunc++] * ydist;
tmpshell += const_wave_f[ifunc++] * zdist;
value += tmpshell * contracted_gto;

contracted_gto = 0.187618 * expf(-0.168714*dist2);
// S_SHELL
value += const_wave_f[ifunc++] * contracted_gto;

contracted_gto = 0.217969 * expf(-0.168714*dist2);
// P_SHELL
tmpshell = const_wave_f[ifunc++] * xdist;
tmpshell += const_wave_f[ifunc++] * ydist;
tmpshell += const_wave_f[ifunc++] * zdist;
value += tmpshell * contracted_gto;

contracted_gto = 3.858403 * expf(-0.800000*dist2);
// D_SHELL
tmpshell = const_wave_f[ifunc++] * xdist2;
tmpshell += const_wave_f[ifunc++] * ydist2;
tmpshell += const_wave_f[ifunc++] * zdist2;
tmpshell += const_wave_f[ifunc++] * xdist * ydist;
tmpshell += const_wave_f[ifunc++] * xdist * zdist;
tmpshell += const_wave_f[ifunc++] * ydist * zdist;
value += tmpshell * contracted_gto;
```

VMD MO JIT Performance Results for C_{60}
2.6GHz Intel X5550 vs. NVIDIA C2050

Kernel	Cores/GPUs	Runtime (s)	Speedup
CPU ICC-SSE	1	30.64	1.0
CPU ICC-SSE	8	4.13	7.4
CUDA-JIT, Zero-copy	1	0.174	176

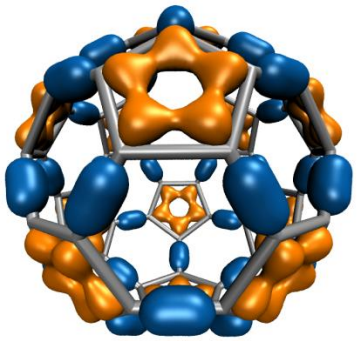
C_{60} basis set 6-31Gd. We used a high resolution MO grid for accurate timings. A more typical calculation has 1/8th the grid points.

JIT kernels eliminate overhead for low trip count for loops, replace dynamic table lookups with constants, and increase floating point arithmetic intensity

Experiments Porting VMD CUDA Kernels to OpenCL

- Why mess with OpenCL?
  - OpenCL is very similar to CUDA, though a few years behind in terms of HPC features, aims to be the "OpenGL" of heterogeneous computing
  - As with CUDA, OpenCL provides a low-level language for writing high performance kernels, until compilers do a much better job of generating this kind of code
  - Potential to eliminate hand-coded SSE for CPU versions of compute intensive code, looks more like C and is easier for non-experts to read than hand-coded SSE or other vendor-specific instruction sets, intrinsics

Molecular Orbital Inner Loop, Hand-Coded SSE Hard to Read, Isn't It? (And this is the "pretty" version!)

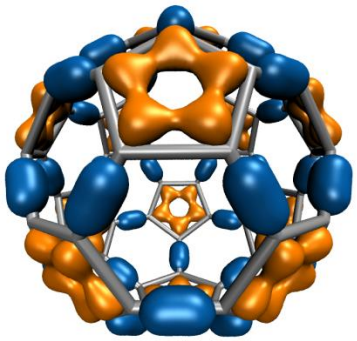


```
for (shell=0; shell < maxshell; shell++) {
  __m128 Cgto = _mm_setzero_ps();
  for (prim=0; prim<num_prim_per_shell[shell_counter]; prim++) {
    float exponent       = -basis_array[prim_counter    ];
    float contract_coeff =  basis_array[prim_counter + 1];
    __m128 expval = _mm_mul_ps(_mm_load_ps1(&exponent), dist2);
    __m128 ctmp = _mm_mul_ps(_mm_load_ps1(&contract_coeff), exp_ps(expval));
    Cgto = _mm_add_ps(contracted_gto, ctmp);
    prim_counter += 2;
  }
  __m128 tshell = _mm_setzero_ps();
  switch (shell_types[shell_counter]) {
    case S_SHELL:
      value = _mm_add_ps(value, _mm_mul_ps(_mm_load_ps1(&wave_f[ifunc++]), Cgto));    break;
    case P_SHELL:
      tshell = _mm_add_ps(tshell, _mm_mul_ps(_mm_load_ps1(&wave_f[ifunc++]), xdist));
      tshell = _mm_add_ps(tshell, _mm_mul_ps(_mm_load_ps1(&wave_f[ifunc++]), ydist));
      tshell = _mm_add_ps(tshell, _mm_mul_ps(_mm_load_ps1(&wave_f[ifunc++]), zdist));
      value = _mm_add_ps(value, _mm_mul_ps(tshell, Cgto));
      break;
```

Until now, writing SSE kernels for CPUs required assembly language, compiler intrinsics, various libraries, or a really smart autovectorizing compiler and lots of luck...

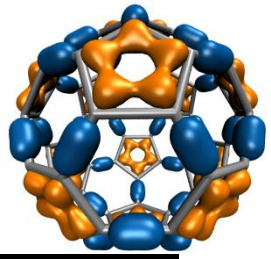
Molecular Orbital Inner Loop, OpenCL Vec4 Ahhh, much easier to read!!!

```
for (shell=0; shell < maxshell; shell++) {
  float4 contracted_gto = 0.0f;
  for (prim=0; prim < const_num_prim_per_shell[shell_counter];  prim++) {
    float exponent       = const_basis_array[prim_counter    ];
    float contract_coeff = const_basis_array[prim_counter + 1];
    contracted_gto += contract_coeff * native_exp2(-exponent*dist2);
    prim_counter += 2;
  }
  float4 tmpshell=0.0f;
  switch (const_shell_symmetry[shell_counter]) {
    case S_SHELL:
      value += const_wave_f[ifunc++] * contracted_gto;      break;
    case P_SHELL:
      tmpshell += const_wave_f[ifunc++] * xdist;
      tmpshell += const_wave_f[ifunc++] * ydist;
      tmpshell += const_wave_f[ifunc++] * zdist;
      value += tmpshell * contracted_gto;
      break;
```



OpenCL's C-like kernel language is easy to read, even 4-way vectorized kernels can look similar to scalar CPU code. All 4-way vectors shown in green.

Apples to Oranges Performance Results: OpenCL Molecular Orbital Kernels



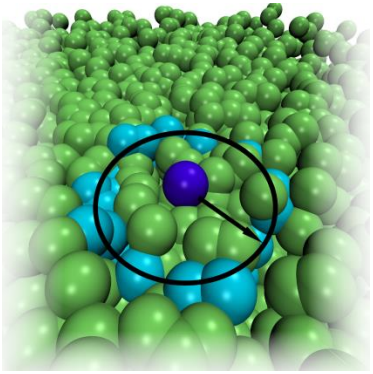
Kernel	Cores	Runtime (s)	Speedup
Intel QX6700 CPU ICC-SSE (SSE intrinsics)	1	46.580	1.00
Intel Core2 Duo CPU OpenCL scalar	2	43.342	1.07
Intel Core2 Duo CPU OpenCL vec4	2	8.499	5.36
Cell OpenCL vec4*** no __constant	16	6.075	7.67
Radeon 4870 OpenCL scalar	10	2.108	22.1
Radeon 4870 OpenCL vec4	10	1.016	45.8
GeForce GTX 285 OpenCL vec4	30	0.364	127.9
GeForce GTX 285 CUDA 2.1 scalar	30	0.361	129.0
GeForce GTX 285 OpenCL scalar	30	0.335	139.0
GeForce GTX 285 CUDA 2.0 scalar	30	0.327	142.4

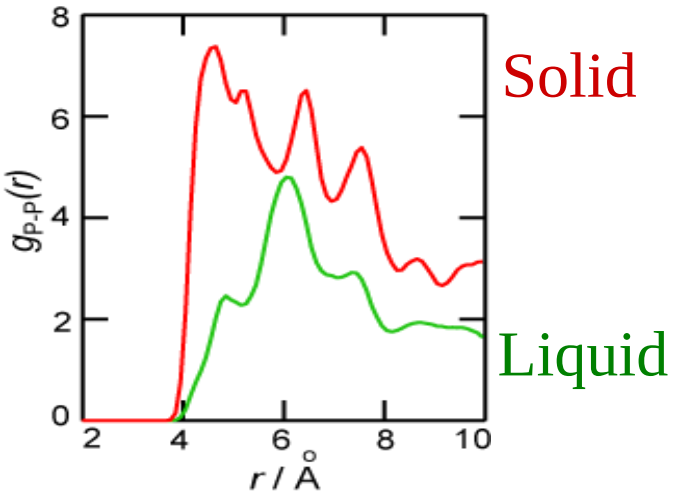
Minor varations in compiler quality can have a strong effect on "tight" kernels. The two results shown for CUDA demonstrate performance variability with compiler revisions, and that with vendor effort, OpenCL has the potential to match the performance of other APIs.


1867

Radial Distribution Function

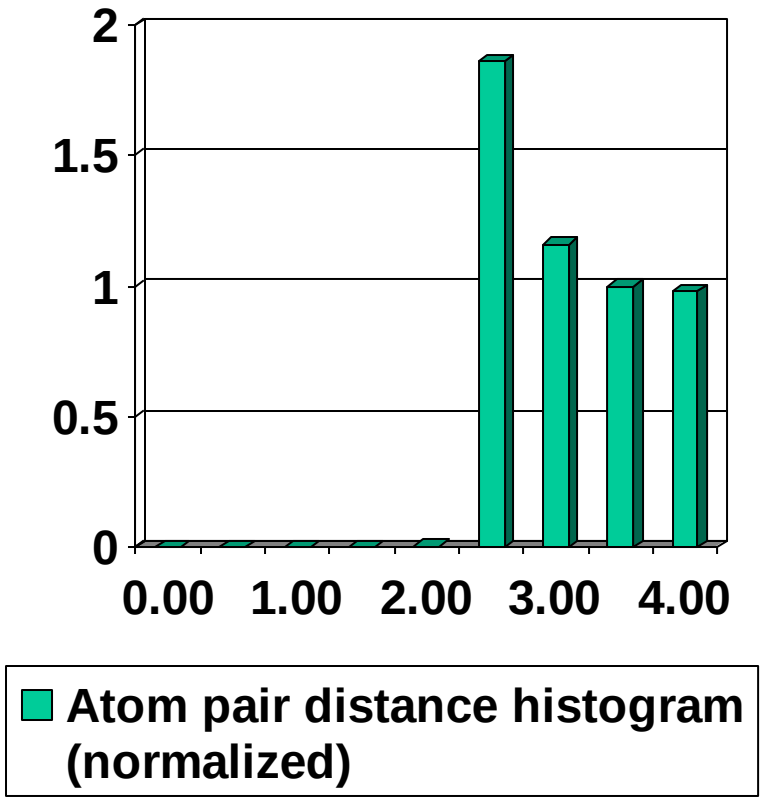
- RDFs describes how atom density varies with distance
- Can be compared with experiments
- Shape indicates phase of matter: sharp peaks appear for solids, smoother for liquids
- Quadratic time complexity $O(N^2)$





Histogramming

- Partition population of data values into discrete bins
- Compute by traversing input population and incrementing bin counters



Computing RDFs

- Compute distances for all pairs of atoms between two groups of atoms A and B
- A and B may be the same, or different
- Use nearest image convention for periodic systems
- Each pair distance is inserted into a histogram
- Histogram is normalized one of several ways depending on use, but usually according to the volume of the spherical shells associated with each histogram bin

Computing RDFs on CPUs

- Atom coordinates can be traversed in a strictly consecutive access pattern, yielding good cache utilization
- Since RDF histograms are usually small to moderate in size, they normally fit entirely in L2 cache
- CPUs can compute the entire histogram in a **single pass**, regardless of the problem size or number of histogram bins

Histogramming on the CPU (slow-and-simple C)

```
memset(histogram, 0, sizeof(histogram));
for (i=0; i<numdata; i++) {
  float val = data[i];
  if (val >= minval  && val <= maxval) {
    int bin = (val - minval) / bindelta;
    histogram[bin]++;
  }
}
```

Fetch-and-increment:

random access updates to histogram bins…

Parallel Histogramming on Multi-core CPUs

- Parallel updates to a single histogram bin creates a **potential output conflict**
- CPUs have atomic increment instructions, but they often take hundreds of clock cycles; unsuitable…
- **SSE can't be used effectively:** lacks ability to "scatter" to memory (e.g. no *scatter-add*, no indexed store instructions)
- For small numbers of CPU cores, it is best to **replicate** and **privatize** the histogram for each CPU thread, compute them independently, and combine the separate histograms in a final reduction step

Computing RDFs on the GPU

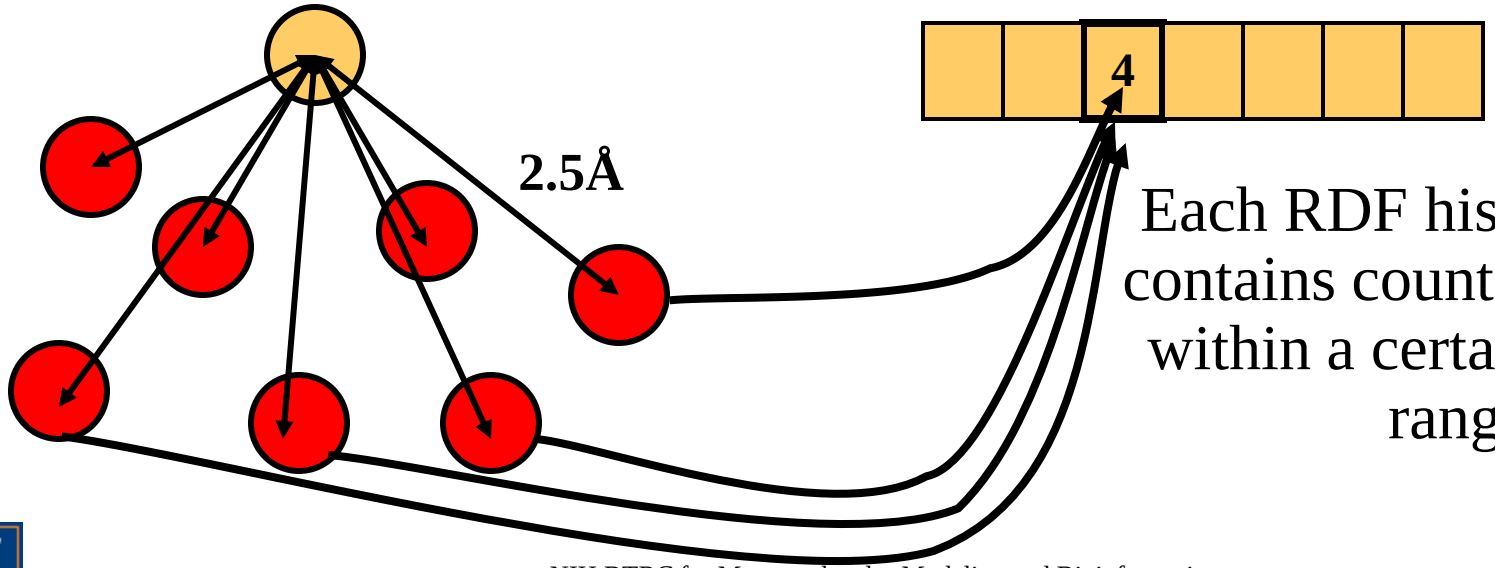
- Need tens of thousands of independent threads
- Each GPU thread computes one or more atom pair distances
- Performance is limited by the speed of histogramming
- Histograms are best stored in fast on-chip shared memory
- Small size of shared memory severely constrains the range of viable histogram update techniques
- **Fast CUDA implementation on Fermi: 30-92x faster than CPU**

Computing Atom Pair Distances on the GPU

- Memory access pattern is simple
- Primary consideration is **amplification of effective memory bandwidth**, through use of GPU on-chip shared memory, caches, and broadcast of data to multiple or all threads in a thread block

Radial Distribution Functions on GPUs

- Load blocks of atoms into shared memory and constant memory, compute periodic boundary conditions and atom-pair distances, all in parallel…
- Each thread computes all pair distances between its atom and all atoms in constant memory, incrementing the appropriate bin counter in the RDF histogram..



Each RDF histogram bin contains count of particles within a certain distance range

GPU Histogramming

- Tens of thousands of threads concurrently computing atom distance pairs…
- Far too many threads for a simple per-thread histogram privatization approach like CPU…
- Viable approach: **per-warp histograms**
- Fixed size shared memory limits histogram size that can be computed in a single pass
- Large histograms require **multiple passes**, but we can skip block pairs that are known not to contribute to a histogram window

Per-warp Histogram Approach

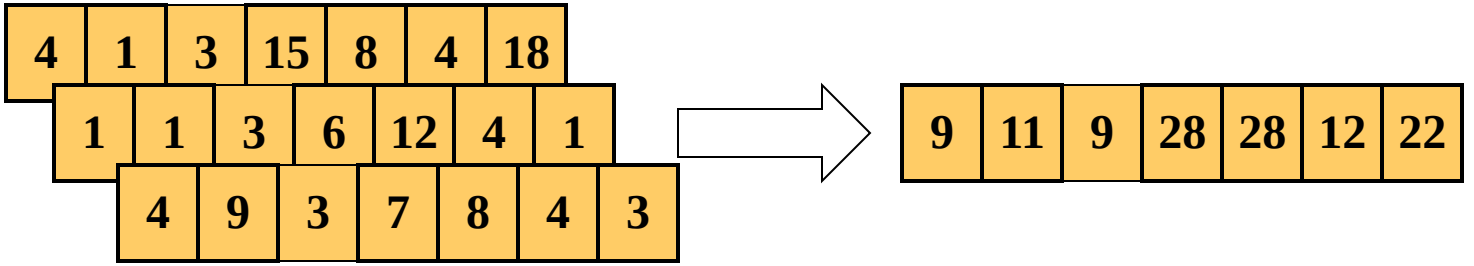
- Each warp maintains its own **private** histogram in on-chip shared memory
- Each thread in the warp computes an atom pair distance and updates a histogram bin in parallel
- Conflicting histogram bin updates are resolved using one of two schemes:
  - Shared memory write combining with thread-tagging technique (older hardware, e.g. G80, G9x)
  - **atomicAdd()** to shared memory (new hardware)

RDF Inner Loops (abbreviated, xdist-only)

```
// loop over all atoms in constant memory
for (iblock=0; iblock<loopmax2; iblock+=3*NCUDABLOCKS*NBLOCK) {
   __syncthreads();
   for (i=0; i<3; i++) xyzi[threadIdx.x + i*NBLOCK]=pxi[iblock + i*NBLOCK]; // load coords…
   __syncthreads();
   for (joffset=0; joffset<loopmax; joffset+=3) {
      rxij=fabsf(xyzi[idxt3  ] - xyzj[joffset  ]);  // compute distance, PBC min image convention
      rxij2=celld.x - rxij;
      rxij=fminf(rxij, rxij2);
      rij=rxij*rxij;
      […other distance components…]
      rij=sqrtf(rij + rxij*rxij);
      ibin=__float2int_rd((rij-rmin)*delr_inv);
      if (ibin<nbins && ibin>=0 && rij>rmin2) {
        atomicAdd(llhists1+ibin, 1U);
      }
   } //joffset
} //iblock
```

Writing/Updating Histogram in Global Memory

- When thread block completes, add independent per-warp histograms together, and write to per-thread-block histogram in global memory
- Final reduction of all per-thread-block histograms stored in global memory

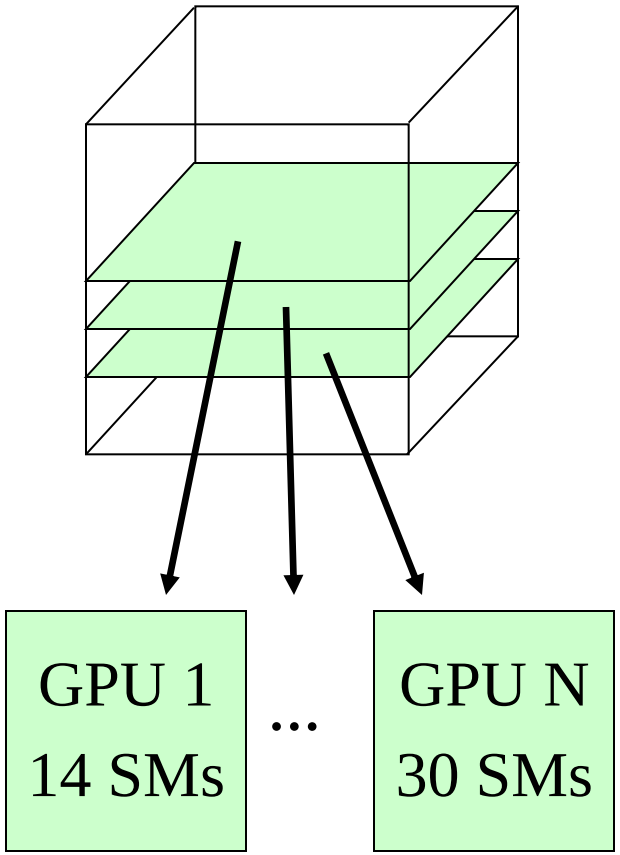


Preventing Integer Overflows

- Since all-pairs RDF calculation computes many billions of pair distances, we have to **prevent integer overflow** for the 32-bit histogram bin counters (supported by the **atomicAdd()** routine)
- We compute full RDF calculation in **multiple kernel launches**, so each kernel launch computes partial histogram
- Host routines read GPUs and increment large (e.g. long, or double) histogram counters in host memory after each kernel completes

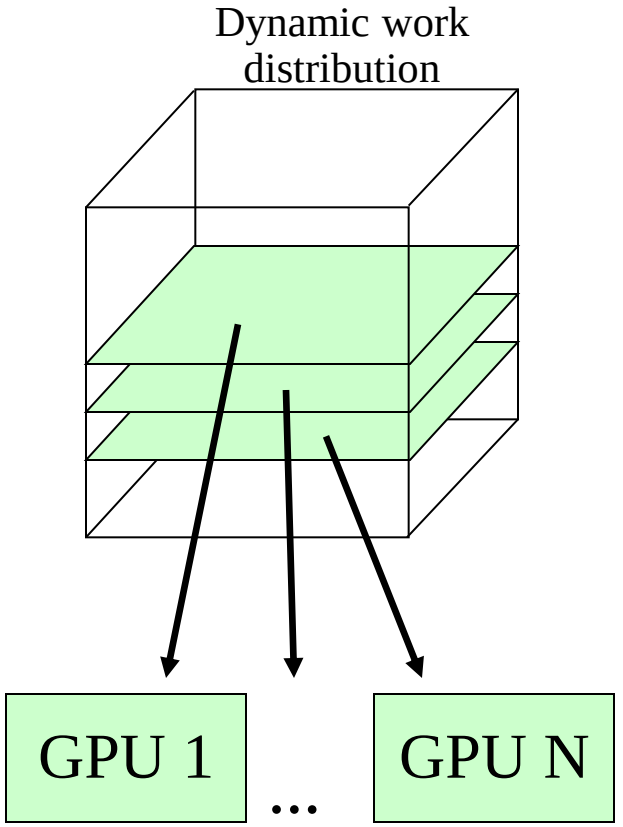
Multi-GPU Load Balance

- Many early CUDA codes assumed all GPUs were identical
- Host machines may contain a diversity of GPUs of varying capability (discrete, IGP, etc)
- Different GPU on-chip and global memory capacities may need different problem “tile” sizes
- Static decomposition works poorly for non-uniform workload, or diverse GPUs



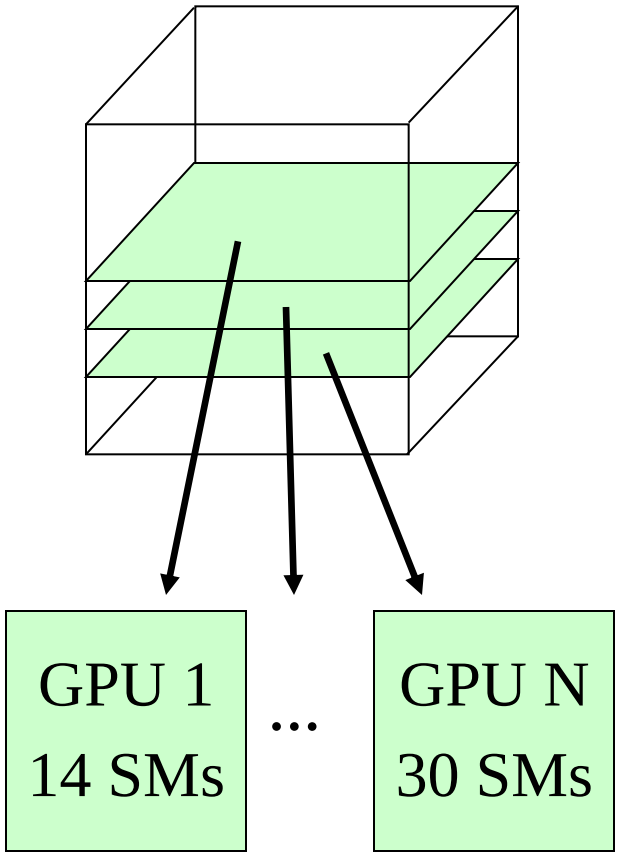
Multi-GPU Dynamic Work Distribution

```
// Each GPU worker thread loops over
// subset of work items…
while (!threadpool_next_tile(&parms,
    tilesize, &tile){
  // Process one work item…
  // Launch one CUDA kernel for each
  //   loop iteration taken…
  // Shared iterator automatically
  //   balances load on GPUs
}
```



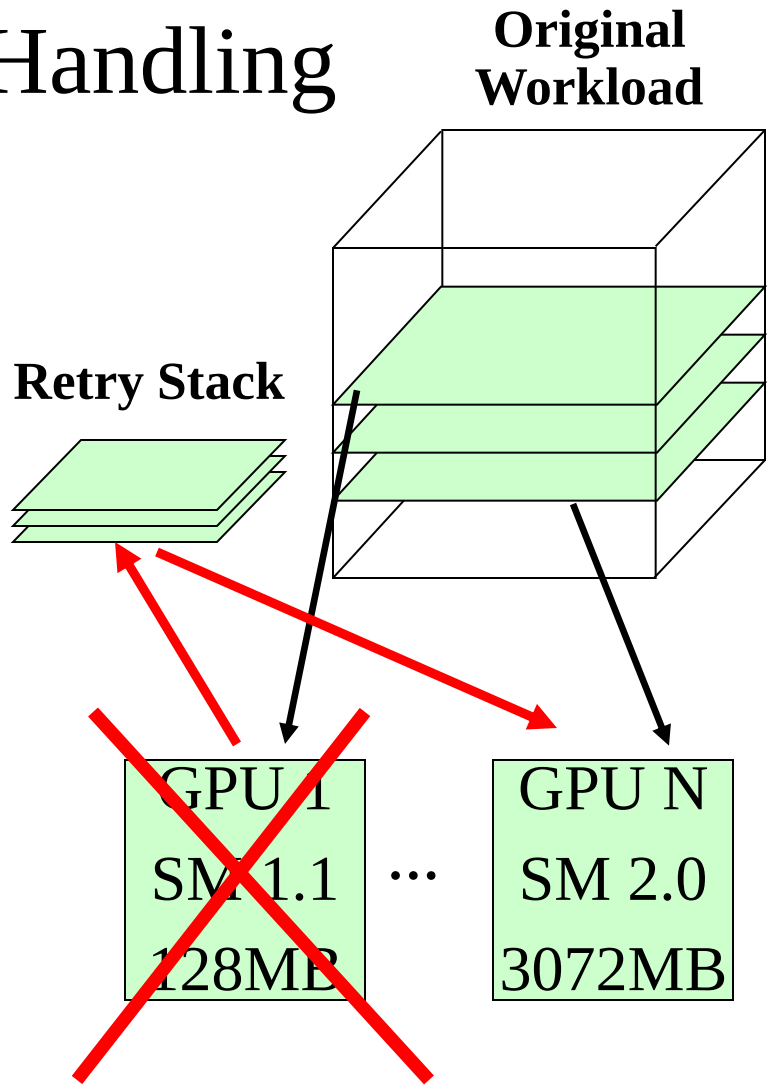
Multi-GPU RDF Calculation

- Distribute combinations of tiles of atoms and histogram regions to different GPUs
- Decomposed over two dimensions to obtain enough work units to balance GPU loads
- Each GPU computes its own histogram, and all results are combined for final histogram



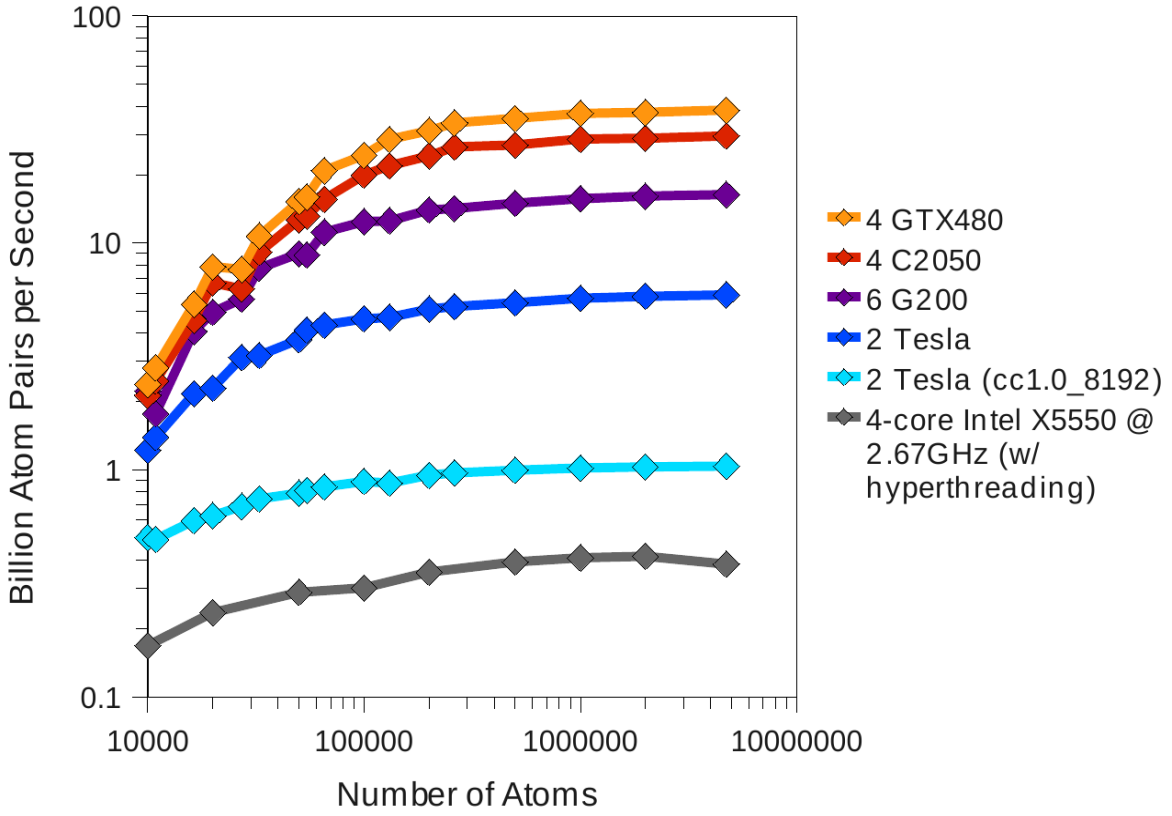
Multi-GPU Runtime Error/Exception Handling

- Competition for resources from other applications can cause runtime failures, e.g. GPU out of memory half way through an algorithm
- Handle exceptions, e.g. convergence failure, NaN result, insufficient compute capability/features
- Handle and/or reschedule failed tiles of work



Multi-GPU RDF Performance

- 4 NVIDIA GTX480 GPUs 30 to 92x faster than 4-core Intel X5550 CPU
- Fermi GPUs ~3x faster than GT200 GPUs: larger on-chip shared memory



Fast Analysis of Molecular Dynamics Trajectories with Graphics Processing Units – Radial Distribution Functions. B. Levine, J. Stone, and A. Kohlmeyer. *J. Comp. Physics*, 230(9):3556-3569, 2011.

Acknowledgements

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GPU Computing Publications

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- **Runtime and Architecture Support for Efficient Data Exchange in Multi-Accelerator Applications** Javier Cabezas, Isaac Gelado, John E. Stone, Nacho Navarro, David B. Kirk, and Wen-mei Hwu. IEEE Transactions on Parallel and Distributed Systems, 2014. (Accepted)
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